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Higham's scaled method for polar decomposition and numerical matrix-inversion

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We present the theory (and illustrating experiments) of the numerical Higham's scaled method for computing the unitary factor of a nonsingular matrix. We show how the quality of the computed inverses of matrices influences the accuracy of the computed polar factorization. In particular: the inversion *via* GECP-factorization and reasonable scaling guarantee a good quality of the computed polar factors (with GEPP-factorization the computed unitary factor can be unacceptable). Some problems of scaling and switching criteria are discussed and experimentally investigated.

Keywords: polar decomposition of a matrix, Higham's method, rounding-errors analysis, numerical matrix-inversion

AMS subject classification: 65G50, 65F30

1 Introduction

We deal with the polar decomposition of a complex nonsingular matrix $A \in \mathbb{C}^{n \times n}$:

$$A = U_A H_A, \quad U_A - \text{unitary}, \quad H_A \in \mathcal{HPD}, \quad (1.1)$$

where \mathcal{HPD} is the class of *Hermitian positive-definite matrices*. U_A is the unitary factor of A (the *orthogonal factor* of $A \in \mathbb{R}^{n \times n}$). Matrices $\{U_A, H_A\}$ are the *polar factors* of A .

The factorization (1.1) can be computed from SVD, the *singular value decomposition* of A . The iterative methods are alternative ways to compute (1.1), see for example [1, 3, 4, 5, 7].

In *Higham's scaled method* [4, 7], denoted by **HS** (referred also as *Newton's scaled method*), one constructs a sequence $\{X_k\}_{k=0}^{\infty}$ of matrices:

$$X_0 = A, \quad X_{k+1} = \frac{1}{2} \left(\gamma_k X_k + \frac{1}{\gamma_k} X_k^{-H} \right), \quad \gamma_k > 0, \quad (1.2)$$

convergent to U_A , the common unitary factor of all X_k . There are known several *theoretical* or *practical* rules of the choice of *scaling parameters* γ_k which increase the speed of convergence, see [3, 4, 7].

Let $\{\tilde{X}_k\}_{k=0}^l$ be the sequence of *iterates* computed in the *numerical HS-algorithm*. In all cases when this algorithm *converges a good unitarity* of the *computed unitary factor* $\tilde{U} \stackrel{\text{df}}{=} \tilde{X}_l$ is achieved:

$$\|\tilde{U}^H \tilde{U} - I\|_2 \leq \varepsilon_0 \quad (1.3)$$

(all ε_s in this paper are of the size ν , the *computing precision*). We can now compute the *Hermitian factor* \tilde{H} of A :

$$\tilde{B} := \tilde{U}^H * A; \quad \tilde{H} := (\tilde{B} + \tilde{B}^H)/2. \quad (1.4)$$

The problem is: *whether* the *computed polar factors* $\{\tilde{U}, \tilde{H}\}$ of A are *acceptable*? That means: *whether* the following *relations* hold:

$$\tilde{H} \in \mathcal{HPD}, \quad \|\tilde{U} \tilde{H} - A\|_2 \leq \varepsilon_1 \|A\|_2 \quad ? \quad (1.5)$$

In [8] we try to explain *how it happens* that the computed by numerical HS polar factors *are acceptable*? We reveal also *two main dangers*: the *poor quality* of the *computed inverses* and using of *too small scaling parameters*.

Our further research is presented in [9]. We explain there all *phenomena* we were able to perceive in *our experiments*. Therefore the *experimental results* play in [9] rather only the role of *illustrations*.

This paper is a *concise version* of [9]. We skip here the proofs, theorem 2.2, the estimation of the accuracy of experimental results and many detailed remarks. We concentrate on the *most important problem* of the *quality* of the *matrix-inversion* in the *numerical HS-algorithm*.

The theory is presented in sections 2 and 4. Section 3 explains *how to read experimental results*. Sections 5 and 6 present briefly some problems of *scaling* and *switching criteria*. For *final conclusions* see section 7.

We add the appendix presenting the proof of the **NC-Property** of the *inversion* by B-method, *via GECP-factorization*, see [2].

2 The theory of HS, the numerical Higham's method

Here and in all next sections HS means the *numerical* HS algorithm (to distinguish from (1.2), where the *theoretical* algorithm is defined).

Let \tilde{X}_k be the *computed iterate* and X_k the *matrix satisfying the conditions* (2.3) below. Neither \tilde{X}_k nor X_k here is identical with X_k in (1.2).

Let us define the following two functions

$$p : \mathbb{C}^{n \times n} \longrightarrow [n^{-1/2}, 1], \quad p(\Psi) \stackrel{\text{df}}{=} \begin{cases} 1, & \text{when } \Psi = 0, \\ \|\Psi\|_2 (\|\Psi\|_F)^{-1}, & \text{otherwise,} \end{cases}$$

$$f : (0, \infty) \longrightarrow [1, \infty), \quad f(t) \stackrel{\text{df}}{=} \frac{1}{2} (t + t^{-1}). \quad (2.1)$$

These *reserved* functions “produce” a series of derivate symbols (f_k, p_k, p_+, \dots) the values of f or p on concrete arguments.

We assume that the computations in HS are performed in the floating-point arithmetic with *precision* ν and that *neither* underflow *nor* overflow occurs.

The *epsilons* ($\varepsilon_0, \varepsilon_x, \dots$) are *modest multiples* of ν . Not all of them must be positive. We signal it writing, for example: $|\varepsilon'_k| \leq \varepsilon$. The only exceptions (see section 4) are “false epsilons” ($\check{\varepsilon}_x, \check{\varepsilon}_k, \dots$), the quantities which ought to be the *true epsilons* (and sometimes are) but – due to breaking of the basic assumption (2.3) – can be much larger than “a modest multiple of ν ”. Usually these false epsilons satisfy $|\check{\varepsilon}| \ll 1$.

Let us formulate already now the following *general assumptions*:

$$\hat{\varepsilon} \text{ cond}_2(A) < 1, \quad \hat{\varepsilon} < \nu^{2/3} \lesssim 10^{-4}, \quad (2.2)$$

where $\hat{\varepsilon}$ is specified in (2.6), (2.3).

2.1 Main definitions and relations

Let us consider a nonsingular matrix $A \in \mathbb{C}^{n \times n}$ and the sequence $\{\tilde{X}_k\}_{k=0}^l$ of matrices (1.2) computed in HS, $\tilde{X}_0 := A$.

Let γ_k be the *chosen scaling parameter* and G_k the *computed inverse* of \tilde{X}_k . We *assume* that *exists* a *nonsingular* matrix X_k *satisfying* the relations:

$$\tilde{X}_k = X_k - \Delta_k, \quad G_k = X_k^{-1} - \Delta'_k, \quad (2.3)$$

where $\|\Delta_k\|_F \leq \varepsilon_x \|X_k\|_2$, $\|\Delta'_k\|_F \leq \varepsilon_g \|X_k^{-1}\|_2$. This defines (not uniquely) X_k for $k < l$. Let us extend it to $k = l$: $X_l = \tilde{X}_l$. The sequences $\{X_k\}$ and $\{\tilde{X}_k\}$ are *neighbour-sequences* and many important properties of X_k are

close to these of \tilde{X}_k . We describe the HS-process in terms of the sequence $\{X_k\}$ since this sequence *imitates well* the relation (1.2), see below (2.5), (2.6).

The *assignment-statements*

$$G_k := \tilde{X}_k^{-1}, \quad \tilde{X}_{k+1} := \left(\tilde{X}_k * \gamma_k + G^H / \gamma_k \right) / 2, \quad (2.4)$$

and (2.3) imply the *equalities*

$$X_{k+1} = Z_{k+1} + T_k, \quad Z_{k+1} \stackrel{\text{df}}{=} \frac{1}{2} \left(\gamma_k X_k + \frac{1}{\gamma_k} X_k^{-H} \right) \quad (2.5)$$

and the *bound*

$$\|T_k\|_F \leq \hat{\varepsilon} f_k, \quad f_k \stackrel{\text{df}}{=} \|Z_{k+1}\|_2, \quad \hat{\varepsilon} = 2\varepsilon_x + \varepsilon_g + 3\sqrt{n}\nu + O(\nu^2). \quad (2.6)$$

Let us consider the SVD of X_k :

$$X_k = P_k \text{diag}(\sigma_1^{(k)}, \dots, \sigma_n^{(k)}) Q_k^H, \quad P_k, Q_k \text{ unitary}$$

and define d_k , the *distance* of X_k from the *unitarity*:

$$d_k \stackrel{\text{df}}{=} \max_i |\sigma_i^{(k)} - 1| = \max \left\{ \sigma_{\max}^{(k)} - 1, 1 - \sigma_{\min}^{(k)} \right\}, \quad (2.7)$$

$$\sigma_{\max}^{(k)} \stackrel{\text{df}}{=} \max_i \{\sigma_i^{(k)}\}, \quad \sigma_{\min}^{(k)} \stackrel{\text{df}}{=} \min_i \{\sigma_i^{(k)}\}. \quad (2.8)$$

The *efficiency* of HS depends on *how quickly* $\{d_k\}_{k=1}^l$ decrease, the *near-unitarity* of the *computed factor* $\tilde{U} = \tilde{X}_l$ depends on the *limiting accuracy* $d \stackrel{\text{df}}{=} \limsup d_k$ of the conceptual infinite sequence $\{d_k\}_{k=0}^\infty$. The *last iterate* \tilde{X}_l constructed in HS should be the *first one* reaching the level $d_l \lesssim d$.

Let us define further quantities

$$c_k \stackrel{\text{df}}{=} \text{cond}_2(X_k) = \frac{\sigma_{\max}^{(k)}}{\sigma_{\min}^{(k)}}, \quad \gamma_k^{(\text{opt})} \stackrel{\text{df}}{=} \left(\sigma_{\max}^{(k)} \sigma_{\min}^{(k)} \right)^{-1/2},$$

$$\rho_k \stackrel{\text{df}}{=} \left(\gamma_k^{(\text{opt})} \gamma_k^{-1} \right)^{-2}, \quad \tau_k \stackrel{\text{df}}{=} \max\{\rho_k, \rho_k^{-1}\}. \quad (2.9)$$

The quantities ρ_k, τ_k “measure” the distance of γ_k from $\gamma_k^{(\text{opt})}$, the *optimal scaling parameter*.

In [9] we show the following relations, see (2.6),

$$f_k = f(\sqrt{c_k \tau_k}), \quad d_{k+1} = (1 - \varepsilon_k^*) f_k - 1, \quad |\varepsilon_k^*| \leq \hat{\varepsilon}, \quad (2.10)$$

$$\hat{\varepsilon} f_k < 1 \quad \text{implies} \quad c_{k+1} \leq (1 + \hat{\varepsilon}) f_k (1 - \hat{\varepsilon} f_k)^{-1}. \quad (2.11)$$

The assumptions (2.2), *practical scaling* ($1, \infty$)-scaling [4] or (F)-scaling [7] and appropriate *switching criteria* in HS guarantee that the sequence $\{f_k\}_{k=0}^{l-1}$ is *strictly decreasing* and the bounds $\hat{\varepsilon} f_k < 1, \tau_k < \sqrt{n}$ hold. We find ultimately in [8] that the bound (1.3) is satisfied with

$$\varepsilon_0 \approx 2d_l \leq \varepsilon' \stackrel{\text{df}}{=} \varepsilon_x + \varepsilon_g + 2\sqrt{n}\nu. \quad (2.12)$$

Remarks 2.1.

- (i) In the case of the standard *double-precision* computations and HS with *practical scaling* in most cases $l \leq 10$ holds.
- (ii) In some *special experiments* (see sections 4 and 5) we modify the *normal* HS-algorithm *introducing* (in a few initial steps only) *either* matrices G_k not satisfying (2.3) *or* scaling parameters γ_k much smaller than $\gamma_k^{(\text{opt})}$. But these modifications *neither* destroy the monotonic decrease of $\{f_k\}$ *nor* influence the final convergence of $\{\tilde{X}_k\}$. Hence the bounds (2.12) and (1.3) remain valid.

We need some *further notions* to discuss the *acceptability* (1.5) of the *computed polar factors* $\{\tilde{U}, \tilde{H}\}$.

Let the abbreviations AUF, APF mean: *approximate unitary factor, approximate polar factors*, respectively.

Let us consider any matrices $X, U \in \mathbb{C}^{n \times n}$, X -nonsingular, U -unitary. If $H_{ux} \stackrel{\text{df}}{=} \frac{1}{2}(U^H X + X^H U) \in \mathcal{HPD}$ then we will say that U is an AUF ($\{U, H_{ux}\}$ are APF) of X with *accuracy* (relative error):

$$\text{acc}(U, X) \stackrel{\text{df}}{=} \frac{\|UH_{ux} - X\|_F}{\|X\|_2}.$$

Let us fix now U as the *unitary factor* of $\tilde{U} = \tilde{X}_l = X_l$. Hence the *polar decomposition* of \tilde{U} is, see (2.7), (2.12),

$$\tilde{U} = UH_u, \quad H_u \in \mathcal{HPD}, \quad \|\tilde{U} - U\|_2 = d_l \lesssim \frac{1}{2}\varepsilon'. \quad (2.13)$$

Let now define for $k = 0, \dots, l$ the following matrices and quantities:

$$H_k \stackrel{\text{df}}{=} \frac{1}{2}(U^H X_k + X_k^H U), \quad \delta_k \stackrel{\text{df}}{=} \|X_k - UH_k\|_F \|X_k\|_2^{-1}. \quad (2.14)$$

Evidently the following implication holds: $H_k \in \mathcal{HPD}$ *implies* $\delta_k = \text{acc}(U, X_k)$. In particular, see (2.13), $H_l = H_u \in \mathcal{HPD}$, $\delta_l = \text{acc}(U, X_l) = 0$.

The following lemma shows that the *properties* of the pair $\{H_0, \delta_0\}$ are *decisive* for the *acceptability* of the *computed polar factors* $\{\tilde{U}, \tilde{H}\}$.

Lemma 2.1. Let introduce the quantities $p_0 \stackrel{\text{df}}{=} p(X_0 - UH_0)$, $\varepsilon_I \stackrel{\text{df}}{=} 2.5\varepsilon_x + \varepsilon_g + \nu m(\sqrt{n})$, where $m(t)$ is a modest polynomial in t (depending on the way of computing \tilde{B} in (1.4)). If $(p_0\delta_0 + \varepsilon_I)\text{cond}_2(A) < 1$ holds and $H_0 \in \mathcal{HPD}$ then the following relations hold:

$$\tilde{H} \in \mathcal{HPD}, \quad \left| \frac{\|A - \tilde{U}\tilde{H}\|_2}{\|A\|_2} - p_0\delta_0 \right| \lesssim \varepsilon_I.$$

Remark 2.2. Lemma 2.1 is valid only when G_0 satisfies (2.3).

Conclusion 2.1. The *computed polar factors* $\{\tilde{U}, \tilde{H}\}$ are *acceptable* **iff** $H_0 \in \mathcal{HPD}$, δ_0 is of the order ν and A is sufficiently well-conditioned, since the following bounds hold: $|p_0\delta_0 - \varepsilon_I| \leq \|A - \tilde{U}\tilde{H}\|_2 \|A\|_2^{-1} \leq p_0\delta_0 + \varepsilon_I$.

In the next subsection we present an *explicit expression* of δ_k in terms of: $\delta_{k+1}, \rho_k, c_k, \hat{\varepsilon}$, see (2.17)-(2.20). This opens a chance for “theoretical transfer” from $\delta_l = 0$ to the *important quantity* δ_0 .

We must be prepared that $\text{acc}(U, \tilde{X}_k) \gtrsim \text{acc}(U, \tilde{X}_{k+1})$ holds since the *rounding errors* in the computation of G_k and \tilde{X}_{k+1} , see (2.4), *can partly spoil* the *information* on \tilde{X}_k transferred to \tilde{X}_{k+1} (hence also to $\tilde{U} = \tilde{X}_l$). The same concerns the *neighbour-sequence* $\{X_k\}_{k=0}^l$: the *relation* $\delta_k \gtrsim \delta_{k+1}$ *can be expected!*

We should recognize *benign rounding errors* in (2.4) – such that δ_k is at most *only slightly larger* than δ_{k+1} – and *dangerous rounding errors* – such that $\delta_k \gg \delta_{k+1}$ *can occur*.

2.2 BIT, the backward-induction theorem

Let us introduce the matrix, see (2.5), $\Psi_k \stackrel{\text{df}}{=} UH_{k+1} - Z_{k+1}$ and the quantities, see (2.6), (2.1), (2.8),

$$\xi_k \stackrel{\text{df}}{=} \|\Psi_k\|_2, \quad \vartheta_k \stackrel{\text{df}}{=} \|\Psi_k\|_F f_k^{-1}, \quad r_k \stackrel{\text{df}}{=} \frac{f_k}{f(\sigma_{\max}^{(k)} \gamma_k)}. \quad (2.15)$$

Theorem 2.1 (BIT). If the relations

$$\xi_k < 1, \quad H_{k+1} \in \mathcal{HPD} \quad (2.16)$$

are satisfied then $\delta_k = \vartheta_k |\chi_k + \kappa_k \zeta_k| r_k$, $\zeta_k \stackrel{\text{df}}{=} (3\sqrt{2} + 2)(2 - \xi_k)^{-1} \xi_k$,

$$c_k \vartheta_k |\mu_k + \lambda_k \zeta_k| r_k < 1 \quad \text{implies} \quad H_k \in \mathcal{HPD},$$

where $\chi_k, \mu_k, \kappa_k, \lambda_k$ are real numbers, *either* all equal zero *or* satisfying inequalities:

$$0 \leq \mu_k < \chi_k \leq 1, \quad |\kappa_k| < 1, \quad |\lambda_k| < 1. \quad (2.17)$$

Remark 2.3. Theorem 2.1 is valid also in cases when the matrices G_k, G_{k+1} are not satisfying (2.3).

Corollary 2.1. The quantity r_k , see (2.15), satisfies the relations

$$r_k = \max \{1, (c_k + \rho_k)(c_k \rho_k + 1)^{-1}\} < \max \{1, \rho_k^{-1}\}. \quad (2.18)$$

If the matrices G_k, G_{k+1} satisfy (2.3) then

$$\xi_k = p'_k |\delta_{k+1}(1 + \varepsilon'_k) + \varepsilon'_k| f_k, \quad p'_k \stackrel{\text{df}}{=} p(\Psi_k), \quad |\varepsilon'_k| \leq \hat{\varepsilon}, \quad (2.19)$$

and – provided (2.16) holds –

$$\delta_k = |\delta_{k+1}(1 + \varepsilon'_k) + \varepsilon'_k| |\chi_k + \kappa_k \zeta_k| r_k, \quad |\varepsilon'_k| \leq \hat{\varepsilon}. \quad (2.20)$$

This allows us to *simplify* the *backward-induction rule*: if $\xi_k \ll 1$ and $H_{k+1} \in \mathcal{HPD}$ holds then

$$\delta_k \approx |\delta_{k+1} + \varepsilon'_k| \chi_k r_k, \quad |\varepsilon'_k| \leq \hat{\varepsilon}, \quad \chi_k \in [0, 1], \quad (2.21)$$

$$c_k(\delta_{k+1} + \hat{\varepsilon})(1 + 7\xi_k)r_k < 1 \quad \text{implies} \quad H_k \in \mathcal{HPD}.$$

Remarks 2.4.

- (i) In *double-precision computations* the approximate equality (2.21) *describes adequately* the behaviour of the sequence $\{\delta_k\}$, since in this case all $\{\xi_k\}$ are very small (the only exception can be ξ_0 when G_0 is not satisfying (2.3), see section 4).
- (ii) With *optimal or practical scaling* the relations $\chi_k r_k \lesssim 1$ can be expected, see section 5. But in the *general case* the *rounding errors* in the computations of \tilde{X}_{k+1} in (2.4) *can be dangerous* when $\rho_k \ll 1$ and $c_k \gg 1$ holds: this implies $r_k \gg 1$ (Theorem 2.2 in [9] shows that χ_k tends to decrease with ρ_k , but we can not expect that always $\chi_k r_k \lesssim 1$ holds, see section 5).

- (iii) *Optimal or practical scaling and inverses* G_k, G_{k+1} *satisfying (2.3) guarantee* $\delta_k \lesssim \delta_{k+1} + \hat{\varepsilon}$. Hence in this case the *rounding errors in both operations* of (2.4) are *benign*.
- (iv) If *any* of the matrices G_k, G_{k+1} is not satisfying (2.3) then the bound $\hat{\varepsilon}$ on $|\varepsilon'_k|$ in corollary 2.1 must be replaced with a much larger quantity: the *rounding errors in the computation of such inverse are dangerous*. We deal with such cases in section 4.

3 Introduction to examples of numerical tests

In sections 4 and 5 we present *examples of numerical tests* illustrating relevant fragments of the theory. All our tests were performed for matrices $A \in \mathbb{R}^{n \times n}$, $6 \leq n \leq 35$, in the **IEEE** standard *double-precision*, $\nu = \nu_d \approx 2.2 \times 10^{-16}$ (with cummulation of “inner products” on standard *extended-precision* variables, $\nu = \nu_e \approx 10^{-19}$).

In most cases we present the computed results with at least *two correct leading decimals*. The results marked with a *star* (*) have probably *only one* correct leading decimal. In results with *exclamation mark* (!) even the first decimal is doubtful.

For each example we present the *matrix* A , the information on *matrix-inversion* and *scaling* in **HS**. We present also the quantity $\tilde{\Delta}_l \stackrel{\text{df}}{=} \|\tilde{U}^T \tilde{U} - I\|_F$, $\tilde{U} = \tilde{X}_l$, and the result of the Cholesky-positivity test of \tilde{H} , see (1.4). Then we present for several iterations, $k = 0, 1, \dots$ some of the computed quantities: $c_k, \rho_k, r_k, e_k^{(L)}, e_k^{(R)}, \hat{\delta}_k$ (eventually also some other auxiliary quantities), where

$$e_k^{(L)} \stackrel{\text{df}}{=} \|I - G_k \tilde{X}_k\|_F w_k^{-1}, \quad e_k^{(R)} \stackrel{\text{df}}{=} \|I - \tilde{X}_k G_k\|_F w_k^{-1}, \quad w_k \stackrel{\text{df}}{=} \|\tilde{X}_k\|_2 \|G_k\|_2, \quad (3.1)$$

$$\hat{\delta}_k \stackrel{\text{df}}{=} \|\tilde{X}_k - U \hat{H}_k\|_F \|\tilde{X}_k\|_2^{-1}, \quad \hat{H}_k \stackrel{\text{df}}{=} \frac{1}{2} (U^T \tilde{X}_k + \tilde{X}_k^T U).$$

Remarks 3.1.

- (i) Let $\tilde{p}_0 \stackrel{\text{df}}{=} p(\tilde{X}_0 - U \hat{H}_0)$. Then $\tilde{p}_0 \hat{\delta}_0$ is a *close approximation* of $\|A - \tilde{U} \hat{H}\|_2 \|A\|_2^{-1}$.
- (ii) $\hat{\delta}_k$ is a *close approximation* of δ_k , see (2.14), provided G_k is satisfying (2.3).

Example 3.1. In table 3.1 we present the computed results of the HSTEST-program (see section 3 in [9]) for the 10×10 matrix $A_1 = \text{tril}(\text{rand}(10))^8 \text{rand}(U)$, see [2], $\Delta_9 = 5.14 \times 10^{-18}$, applying (F)-scaling and the GEPP-matrix-inversion. Matrix \tilde{H} passed the positivity test.

Table 3.1

k	$c_k - 1$	ρ_k	$e_k^{(L)}$	$e_k^{(R)}$	$\hat{\delta}_k$
0	$8.74e + 14^*$	0.930^*	$3.10e - 17$	$8.72e - 09$	$5.12e - 09$
1	$1.66e + 06$	0.708	$3.28e - 17$	$1.96e - 15$	$1.19e - 15$
2	$7.56e + 02$	1.00	$5.90e - 17$	$7.52e - 16$	$4.09e - 16$
3	$1.19e + 01$	0.732	$1.07e - 16$	$1.44e - 16$	$2.68e - 16$
4	$1.17e + 00$	1.07	$2.97e - 16$	$2.95e - 16$	$2.80e - 16$
5	$8.38e - 02$	1.03	$5.08e - 16$	$5.16e - 16$	$3.43e - 16$
6	$1.51e - 03$	1.00	$5.74e - 16$	$5.74e - 16$	$3.40e - 16$
7	$7.01e - 07$	1.00	$5.35e - 16$	$5.35e - 16$	$2.64e - 16$
8	$2.46e - 13$	1.00	$4.84e - 16$	$4.84e - 16$	$1.80e - 16$

Remarks 3.2.

- (i) The value of $e_0^{(R)}$ shows that matrix G_0 is not satisfying (2.3).
- (ii) The quantity $\|A - \tilde{U}\tilde{H}\|_2 \|A\|_2^{-1}$, see (1.4) and remark 3.1 (i), cannot be smaller than $\hat{\delta}_0 n^{-1/2} \approx 1.62 \times 10^{-9}$. Hence the computed polar factors $\{\tilde{U}, \tilde{H}\}$ are *not acceptable*. It is the result of breaking the assumption (2.3) for $k = 0$, see section 4.
- (iii) The results presented in table 3.1 for $k > 3$ are typical for all our experiments. In next examples we will present only the relevant part of experimental results.

4 The quality problem of the matrix-inversion in the HS-process

Some contemporary standard procedures *compute* the *inverses* from the *Gaussian triangular factorization with partial pivoting* (GEPP) of the *inverted matrix*, see [2]. Using *these procedures* in the HS-process *yields frequently* (but *not always!*) *acceptable results* (see example 3.1). The inversion *via triangular factorization with complete pivoting* (GECP) yields *practically always acceptable results* in HS with *practical or optimal scaling*.

We should *recognize* the *properties* of the *computed inverse* G_k of \tilde{X}_k *not impeding* the *good numerical behaviour* of the HS-process and *those properties* which *can seriously spoil* the *quality* of the *computed unitary factor* \tilde{U} of A .

4.1 Properties of computed inverses

Let G be the computed inverse of the nonsingular matrix X . We introduce auxiliary quantities $x \stackrel{\text{df}}{=} \|X\|_2$, $g \stackrel{\text{df}}{=} \|G\|_2$, $c \stackrel{\text{df}}{=} \text{cond}_2(X) = x\|X^{-1}\|_2$ and consider the following four eventual properties of G :

$$\|G - X^{-1}\|_F \leq \varepsilon g c, \quad (4.1)$$

$$\|GX - I\|_F \leq \varepsilon g x, \quad (4.2)$$

$$\|XG - I\|_F \leq \varepsilon g x, \quad (4.3)$$

$$\exists \Delta', \Delta : \quad G + \Delta' = (X + \Delta)^{-1}, \quad \|\Delta'\|_F \leq \varepsilon_g g, \quad \|\Delta\|_F \leq \varepsilon_x x. \quad (4.4)$$

The same relations define the *properties of inversion procedures* as follows: Let \mathbb{M} be a subset of nonsingular $n \times n$ matrices X . We say that an inversion algorithm Inv is *numerically stable* (NS) in \mathbb{M} if for each $X \in \mathbb{M}$ the computed inverse G satisfies (4.1). In the same way:

- (4.2) defines the *left-residual stability* (LRS) of Inv in \mathbb{M} ,
- (4.3) defines the *right-residual stability* (RRS) of Inv in \mathbb{M} ,
- (4.4) defines the *numerical correctness* (NC) of Inv in \mathbb{M} .

We shall use the same notation: NS, LRS, RRS, NC for the properties (4.1)–(4.4) of the matrix G (no matter what is the “official property” in \mathbb{M} of the algorithm which computed G).

We define also two *combined* properties of G :

$$\text{Alt} \stackrel{\text{df}}{=} \text{LRS or RRS}, \quad \text{Conj} \stackrel{\text{df}}{=} \text{LRS and RRS}. \quad (4.5)$$

Assuming $\varepsilon_x + \varepsilon_g + \varepsilon_x \varepsilon_g \leq \varepsilon$ and $\varepsilon x g < 1$ we find the following implications:

$$\text{NC} \implies \text{Conj} \implies \text{Alt} \implies \text{NS} \quad (4.6)$$

and the bounds

$$\|GX - I\|_F \leq c \|XG - I\|_F, \quad \|XG - I\|_F \leq c \|GX - I\|_F. \quad (4.7)$$

Let us note further that for small c , say $c \leq 10$, NS implies NC (for example: with $\varepsilon_x = 0, \varepsilon_g \leq 10\varepsilon$). Hence all listed properties of G can differ distinctly only when $c = \text{cond}_2(X)$ is large.

Further definitions: We will say that G has LRS-Only-Property if G has the LRS-Property but has not the RRS-Property. In this case G has the Alt-Property (hence also NS-Property) but has *neither* the Conj-Property *nor* the NC-Property. In the same way, using the term: to have the Only-Property, we define other eventual *highest-properties* of G in the hierarchical system defined by (4.5), (4.6).

Let us note at last that the NC-Property is the *highest general quality* (expressed in *norms of matrices*) of an *inverse G computed in a constant finite precision*. According to the formulation of W. Kahan, see [6], in this NC-case: G is a *slightly wrong inverse* of a *slightly wrong matrix X* .

4.2 The W-conjecture

There are several versions of computing the inverse G from GEPP-triangular factorization of X , see [2], which are either *left-residual* - or *right-residual-stable* in a broad subset \mathbb{M} of $n \times n$ matrices. Hence such *GEPP-inversion algorithms* guarantee the Alt-Property of computed inverses. For well-conditioned matrices X it means practically the Conj-Property of G . But also for badly conditioned matrices X we can check directly that *frequently both residuals* $\|GX - I\|_F, \|XG - I\|_F$ *are small* (are bounded by εxg), see [2], [10]. That means that G has the Conj-Property in spite of (4.7) with large c .

J.H. Wilkinson *explained this phenomenon*, in [10, pp. 110-111], showing that the matrix G (computed *via GEPP-factorization by A-method*, see [2]) has the NC-Property provided the *triangular systems* – involved in the computation of G from GEPP – are *solved to high accuracy*. This happens *frequently but not always*. It seems probable that *this is the only reason why happens* the Conj-Property of inverses computed *via GEPP-factorization*. Let us express it as follows:

W-conjecture. If an inverse G computed *via GEPP-factorization* of X has the Conj-Property then, probably, G has also the (stronger) NC-Property.

The experiments of subsection 4.5 and all our experiments with GEPP-inversion seem to justify the W-conjecture.

4.3 HS with inverses not always satisfying (2.3)

In (2.3) we postulate in fact the NC-property (4.4) of all computed inverses G_k of $\tilde{X}_k, k = 0, \dots, l-1$. Hence, see remark 2.4 (iii), the NC-property of all $\{G_k\}$ is *sufficient* for *good behaviour* of the HS-process with *practical scaling*.

The problem is *whether* the *inverses* G_k *not satisfying* (2.3) *can spoil* (and *how much?*) the *quality* of the *computed unitary factor* \tilde{U} ?

We will consider only the case of G_k with **Alt-Property** (this includes the **Conj-Property** and **NC-Property** as special subcases). Let us incorporate *these eventual deviations* (from the *normality* of (2.3)) into our general description of HS.

Let assume hence the relations

$$X_k = \tilde{X}_k + \Delta_k, \quad X_k^{-1} = G_k + \Delta'_k \quad (4.8)$$

and let us introduce the quantities (in general: *false epsilons*): $\tilde{\varepsilon}_x^{(k)} \stackrel{\text{df}}{=} \|\Delta_k\|_F \|X_k\|_2^{-1}$, $\tilde{\varepsilon}_g^{(k)} \stackrel{\text{df}}{=} \|\Delta'_k\|_F \|X_k^{-1}\|_2^{-1}$.

Let us assume further the relations: $\hat{c}_k \gg 1$, $\varepsilon \hat{c}_k \ll 1$, $\hat{c}_k \stackrel{\text{df}}{=} \text{cond}_2(\tilde{X}_k)$. We present below a *simplified version* of theorem 4.1 in [9], using an *approximate equality* $a \approx b$ (a, b nonnegative) meaning any of the following *three possibilities*:

$$|a - b| \leq O(\varepsilon), \quad |a - b| \leq O(\varepsilon \hat{c}_k) \max\{a, b\}, \quad |a - b| \leq O(\hat{c}_k^{-1}) \max\{a, b\}.$$

Theorem 4.1. The only minimizer Δ_k of the linear functional

$$\varphi_k(\Delta) \stackrel{\text{df}}{=} \max_{\Delta \in \mathbb{C}^{n \times n}} \left\{ \frac{\|\Delta\|_F}{\|\tilde{X}_k\|_2}, \frac{\|\tilde{X}_k^{-1} - G_k - \tilde{X}_k^{-1} \Delta \tilde{X}_k^{-1}\|_F}{\|G_k\|_2} \right\}$$

defines in (4.8) the nonsingular matrix X_k and the matrix Δ'_k such that the following relations hold:

$$c_k \stackrel{\text{df}}{=} \text{cond}_2(X_k) \approx \hat{c}_k, \quad \tilde{\varepsilon}_x^{(k)} \approx \tilde{\varepsilon}_g^{(k)} \approx \hat{\varphi}_k \stackrel{\text{df}}{=} \varphi_k(\Delta_k).$$

Introducing the quantities, see (3.1):

$$\varepsilon_k^{(A)} \stackrel{\text{df}}{=} \min \left\{ e_k^{(L)}, e_k^{(R)} \right\}, \quad \tilde{\varepsilon}_k^{(A)} \stackrel{\text{df}}{=} \max \left\{ e_k^{(L)}, e_k^{(R)} \right\}, \quad e_k^{(c)} = \sqrt{e_k^{(L)} e_k^{(R)}},$$

we specify $\hat{\varphi}_k$ according to the assumed property of G_k :

(i) If G_k has the **Alt-Only-Property** then $\hat{\varphi}_k = \hat{\varphi}_k^{(\text{Alt})}$ where

$$\frac{1}{2}\check{\varepsilon}_k^{(\text{A})} \lesssim \hat{\varphi}_k^{(\text{Alt})} \lesssim \frac{1}{\sqrt{2}}e_k^{(c)}c_k^{1/2} \lesssim \frac{1}{\sqrt{2}}\varepsilon_k^{(\text{A})}c_k. \quad (4.9)$$

(ii) If G_k has the **Conj-Only-Property** then $\hat{\varphi}_k = \hat{\varphi}_k^{(\text{Conj})}$ where

$$\frac{1}{2}\varepsilon_k^{(\text{A})} \lesssim \hat{\varphi}_k^{(\text{Conj})} \lesssim \frac{1}{\sqrt{2}}\varepsilon_k^{(c)}c_k^{1/2}, \quad \varepsilon_k^{(c)} \stackrel{\text{df}}{=} e_k^{(c)}. \quad (4.10)$$

(iii) If G_k has the **NC-Property** then $\hat{\varphi}_k = \hat{\varphi}_k^{(\text{NC})}$ where $\hat{\varphi}_k^{(\text{NC})} \lesssim \max\{\varepsilon_x, \varepsilon_g\}$.

If matrices G_k, G_{k+1} can have the **Alt** – or **Conj** – or **NC-Property** then the bound $\hat{\varepsilon}$ in the relevant relations of section 2 must be replaced with $\check{\varepsilon}_k^*$: $\check{\varepsilon}_k^* \stackrel{\text{df}}{=} 2\hat{\varphi}_k + \hat{\varphi}_{k+1} + 3\sqrt{n}\nu$. But for important recursive formulas (2.20), (2.21) we should rather choose the *presentation exposing the potentially dominating terms*. For example, when $\xi_k \ll 1, H_{k+1} \in \mathcal{HPD}$ and G_k has the **Alt-Only-Property** or the **Conj-Only-Property**, let us choose the presentation:

$$\delta_k \approx \left| |\varphi_k^* + \theta'_k \hat{\varphi}_{k+1}| + \theta''_k |\delta_{k+1} + O(\nu)| \right| \chi_k r_k, \quad \theta'_k, \theta''_k \in [-1, 1],$$

where $\varphi_k^* \stackrel{\text{df}}{=} \|\Delta_k \gamma_k + \Delta_k'^H \gamma_k^{-1}\|_F (2f_k)^{-1}$. Closer examination of the matrices Δ_k, Δ_k' shows that the following bounds (respectively) hold:

$$\hat{\varphi}_k^{(\text{Alt})} \lesssim \varphi_k^* \lesssim \sqrt{2}\hat{\varphi}_k^{(\text{Alt})} \quad \text{or} \quad \varphi_k^* \lesssim 2\hat{\varphi}_k^{(\text{Conj})}. \quad (4.11)$$

Hence in the case of *distinctly Alt-Only-Property* of G_k (when $\check{\varepsilon}_k^{(\text{A})} \gg \max\{\hat{\varphi}_{k+1}, \delta_{k+1}\}$) the relation

$$\delta_k \gtrsim \frac{1}{2}\check{\varepsilon}_k^{(\text{A})} = \frac{1}{2} \max\{e_k^{(\text{L})}, e_k^{(\text{R})}\} \quad (4.12)$$

is *inevitable*, see (4.9).

Conclusion 4.1. The *rounding errors* in the *computation* of G_k with **Alt-Only-** or **Conj-Only-Property** are *dangerous*.

4.4 Experiments with inverses G_k having the Alt-Only-Property

We apply here in HS the *practical-scaling* and the computation of the inverses G_k via GEPP-factorization of \tilde{X}_k (versions LRS - or RRS-stable).

In example 3.1 we presented already such experiment with the 10×10 matrix $A_1 = \text{tril}(\text{rand}(10))^8 \text{rand}(U)$, see [2].

Example 4.1. The matrices \tilde{H} passed the positivity test only in examples (i), (ii) below. Matrices A_3, A_4 are defined in [2].

- (i) The results for $A_2 = A_1^T, \tilde{\Delta}_9 = 6.2 \times 10^{-16}$ are presented in table 4.1.

Table 4.1

k	c_k	$e_k^{(L)}$	$e_k^{(R)}$	$\hat{\delta}_k$
0	$8.75e + 14^*$	$8.79e - 09$	$3.25e - 17^*$	$5.45e - 09$
1	$1.86e + 06$	$5.57e - 15$	$6.12e - 17^*$	$2.69e - 15$
2	$2.96e + 02$	$6.39e - 16$	$3.46e - 16$	$3.46e - 16$

- (ii) Table 4.2 includes the results for $n = 15, A_3 = \text{rand}(Q)\text{qr}(\text{vand}(15)), \tilde{\Delta}_{10} = 9.17 \times 10^{-16}$.

Table 4.2

k	c_k	$e_k^{(L)}$	$e_k^{(R)}$	$\hat{\delta}_k$
0	$1.58e + 13$	$3.68e - 17^*$	$3.91e - 14$	$2.13e - 14$
1	$1.11e + 06$	$8.92e - 17^*$	$1.65e - 14$	$8.23e - 15$
2	$4.82e + 02$	$1.38e - 16$	$1.21e - 15$	$7.12e - 16$
3	$1.15e + 01$	$2.22e - 16$	$3.01e - 16$	$5.47e - 16$

- (iii) In table 4.3 we give the results for $n = 25, A_4 = \text{rand}(Q)\text{qr}(\text{vand}(25)), \tilde{\Delta}_{10} = 2.46 \times 10^{-15}$.

Table 4.3

k	c_k	$e_k^{(L)}$	$e_k^{(R)}$	$\hat{\delta}_k$
0	$1.87e + 18!$	$2.93e - 17^*$	$1.39e - 10$	$8.55e - 11$
1	$4.25e + 08$	$8.65e - 17^*$	$1.67e - 12$	$7.67e - 13$
2	$1.10e + 04$	$1.15e - 16$	$6.69e - 15$	$3.75e - 15$
3	$5.26e + 01$	$3.47e - 16$	$6.38e - 16$	$1.09e - 15$

Remarks 4.1.

- (i) For example (i) see remarks 3.1 (i), (ii).
- (ii) In examples (ii), (iii) G_0 and G_1 have the **LRS-Only-Property**.
- (iii) Notice that the relation (4.12) is clearly demonstrated for $k = 0$ in all tests of example 3.1 and example 4.1.

4.5 Experiments with inverses G_k having the Conj-Only-Property

We apply here in HS the *optimal-scaling* and the procedure `INVCONJ(X)` yielding (via SVD of X) the computed inverse G of X with **Conj-Property** (if possible: with **Conj-Only-Property**), see subsection 4.5 in [9]. We present below the experiments with matrices $A_s = P_s \text{diag}(\sigma_j^{(s)}) Q_s^T \in \mathbb{R}^{n \times n}$ for $s = 5, 6, 7$ (P_s, Q_s orthogonal, random). In all these experiments the relative residuals $e_k^{(L)}, e_k^{(R)}$ are not exceeding 2.7×10^{-15} . Hence we present only the quantities $c_k, c_k^{1/2}, \hat{\delta}_k, m_k$, where m_k is the number of singular values $\{\hat{\sigma}_i^{(k)}\}$ of X_k close to $\hat{\alpha}_k \stackrel{\text{df}}{=} \left(\hat{\sigma}_{\max}^{(k)} \hat{\sigma}_{\min}^{(k)} \right)^{1/2}$ (with $A_s \in \mathbb{R}^{n \times n}$ the *rounding errors* in G_k with **Conj-Only-Property** are *dangerous* only when $m_k \geq 2$ holds, see subsection 4.5 in [9]).

Examples 4.2. In experiments below all matrices \tilde{H} passed the positivity test.

- (i) In table 4.4 we present the results for $n = 6, \tilde{\Delta}_6 = 5.76 \times 10^{-16}$ and

$$\{\sigma_i^{(5)}\} = \{10^7, \sqrt{2 \times 10^7}, 1, 1, \sqrt{5 \times 10^{-8}}, 10^{-7}\}.$$

Table 4.4

k	c_k	$\sqrt{c_k}$	$\hat{\delta}_k$	m_k
0	1.00e + 14	1.00e + 07	5.49e - 10	2
1	5.06e + 06	2.25e + 03	1.01e - 13	2
2	1.06e + 03	3.26e + 01	8.74e - 16	—

- (ii) In table 4.5 we present the results for $n = 20, \tilde{\Delta}_6 = 1.99 \times 10^{-15}$ and

$$\{\sigma_i^{(6)}\} = \{10^{14}, 10^7, \dots, 10^7, 1\}.$$

Table 4.5

k	c_k	$\sqrt{c_k}$	$\hat{\delta}_k$	m_k
0	$9.99e + 13$	$1.00e + 07$	$7.04e - 09$	18
1	$5.17e + 06$	$2.27e + 03$	$1.72e - 15$	—

(iii) In table 4.6 we present the results for $n = 20$, $\tilde{\Delta}_8 = 1.87 \times 10^{-15}$ and

$$\sigma_i^{(7)} = (10^{14/19})^{i-1} \quad (i = 1, \dots, 20).$$

Table 4.6

k	c_k	$\sqrt{c_k}$	$\hat{\delta}_k$	m_k
0	$1.00e + 14$	$1.00e + 07$	$4.39e - 10$	2
1	$3.61e + 06$	$1.90e + 03$	$1.31e - 13$	2
2	$7.27e + 02$	$8.50e + 01$	$6.62e - 15$	1
3	$1.35e + 01$	$3.07e + 00$	$2.10e - 15$	—

Remark 4.2. The experimental results presented above are evidently consistent with the bounds (4.11), (4.10).

5 The problems of scaling

Assuming: $\xi_k \ll 1$, $H_{k+1} \in \mathcal{HPD}$ and G_k, G_{k+1} satisfying (2.3), we can use the simplified form of recursion, see (2.21),

$$\delta_k \approx |\delta_{k+1} + \varepsilon'_k| z_k, \quad z_k \stackrel{\text{df}}{=} \chi_k r_k, \quad |\varepsilon'_k| \leq \hat{\varepsilon},$$

where: $\chi_k \leq 1$, $r_k = \max\{1, (c_k + \rho_k)(c_k \rho_k + 1)^{-1}\}$, $\rho_k = (\gamma_k / \gamma_k^{(\text{opt})})^2$. If $\gamma_k \ll \gamma_k^{(\text{opt})}$ and $c_k \gg 1$ then $r_k \gg 1$ holds. Though χ_k tends to decrease with ρ_k , see theorem 2.2 in [9], it can happen that also $z_k \gg 1$ holds, what implies $\delta_k \gg \delta_{k+1}$. That is *the problem of too small scaling parameters*.

This can happen in *one step*, but also in *several consecutive steps*, when $z_k > 1, z_{k-1} > 1, \dots$ holds.

In HS with *practical scaling* $r_k < \sqrt{n}$ holds, hence the *danger* is *not very serious*. What's more: all *known experiments* seem to indicate that HS *with practical-scaling* is *immune* to the *danger* of too small scaling parameters: the relation $z_k \stackrel{\text{df}}{=} \chi_k r_k \lesssim 1$ is *always observed*. Section 5 in [9] proposes an *explanation* for this phenomenon.

But for *drastically small scaling parameters* the *danger* of $z_k \gg 1$ *really exists!*

Example 5.1. For a random 10×10 matrix A_8 we apply the HS-process with GECP matrix-inversion and – essentially – (F)-scaling, introducing “artificially” very small γ_k for $k = 0, 2, 4$. The results are presented in table 5.1. We additionally compute the quantities $\hat{\chi}_k \stackrel{\text{df}}{=} \hat{\delta}_k r_k^{-1} (\hat{\delta}_{k+1} + 10^{-16})^{-1}$ (probably lower bounds on χ_k). Matrix \tilde{H} passed the positivity test.

Table 5.1

k	c_k	ρ_k	r_k	$\hat{\delta}_k$	$\hat{\chi}_k$
0	$9.61e + 14^*$	$8.21e - 05^*$	$1.21e + 04^*$	$3.96e - 13$	0.0013
1	$1.12e + 09$	$1.12e + 00$	1	$2.46e - 14$	0.422
2	$1.17e + 04$	$1.27e - 04$	$5.13e + 03$	$5.85e - 14$	0.013
3	$5.17e + 03$	$1.08e + 00$	1	$6.71e - 16$	0.647
4	$3.15e + 01$	$3.25e - 02$	$1.55e + 01$	$8.36e - 16$	0.154
5	$1.64e + 01$	$1.37e + 00$	1	$1.51e - 16$	0.302

Remarks 5.1.

- (i) Table 5.1 demonstrates the tendency of χ_k to decrease with ρ_k .
- (ii) Very small ρ_k (hence large τ_k , (2.9)) retard the decreasing of $\{c_k\}$, see relations (2.10), (2.11).
- (iii) Section 5 in [9] presents more examples of this type.

Another problem is the influence of scaling on the effectiveness of the HS-process. Both considered above ways of *practical scaling* have two advantages:

- for large n , say $n \geq 10$, the cost of computing of $\{\gamma_k\}$ is negligible (with respect to the cost of the matrix-inversion),
- there is a chance of accelerating the convergence when there are large gaps in the spectrum of singular values of A .

The following way of quasi-optimal scaling:

- choose positive quantities a_0, b_0 such that $a_0 < \sigma_j(A) < b_0$ holds,
- compute: $\mu_0 := b_0/a_0$, $\gamma_0^{(q)} := (a_0\sqrt{\mu_0})^{-1}$, and for $k > 0$

$$\mu_k := (\mu_{k-1}^{1/2} + \mu_{k-1}^{-1/2})/2, \quad \gamma_k^{(q)} := \mu_k^{-1/2},$$

guarantees the first advantage for all n ; however, it does not have the second advantage.

6 The switching criteria in HS

In our experiments, aimed to study the problems of sections 4 and 5, we tested additionally the *criteria* (proposed in [4], [7], [8]) for *accepting* the *last computed iterate* as the *computed unitary factor* \tilde{U} . We tested also the *criteria* (proposed in [4], [8]) for *switching* from $(1, \infty)$ -*scaling* to *unscaled iterations*. Section 6 in [9] presents the details of these tests. One of the *conclusions* is presented in section 7 (iii).

7 Final conclusions

- (i) Matrix-inversion in the HS-process should yield the computed inverse G of the matrix X (the inverse of the current iterate) satisfying the condition (2.3) (the NC-property). This property is warranted by the inversion *via* GECP-triangularization of X . Using in HS the standard inversion *via* GEPP, see [2], can fail, yielding for some special matrices A a poor unitary factor \tilde{U} . This will never occur for well-conditioned matrices A , say: $\text{cond}_2(A) \leq 10^2$.
- (ii) Using in the HS-process a good matrix-inversion, see (i), and *either* (F)-scaling [7] *or* $(1, \infty)$ -scaling [4] (with appropriate switch to unscaled iterations) practically guarantees good quality of the computed unitary factor \tilde{U} of A (the same quality, as yields the unitary factor computed *via* SVD of A).
- (iii) An appropriate stopping criterion in most cases guarantees that $\tilde{U} = \tilde{X}_l$ is the first iterate reaching the limiting accuracy. With the stopping criterion in [4] frequently one redundant step is performed.
- (iv) The formal cost (the number of arithmetic operations) of the HS-process in the standard-double precision is at most of the same order as for SVD (is smaller for well-conditioned matrices or matrices with large gap in the spectrum of the singular values).
- (v) Using in the HS-process scaling parameters $\{\gamma_k\}$ distinctly larger or smaller than the optimal ones, see relations (2.9) and (2.10), can spoil the convergence. Using $\{\gamma_k\}$ distinctly smaller is spoiling also the quality of \tilde{U} as an approximate unitary factor of A . Practical scaling, see (iii), is not involving such impendency.

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A1 Numerical correctness of inverting matrices via GECP

Theorem A1.1. Let assume that the GECP-process for $n \times n$ matrix X yields the permutation matrices P_L, P_R , and lower and upper triangular matrices $L = \text{tril}(L) = [l_{ij}], R = \text{triu}(R) = [r_{ij}]$, respectively, such that the following relations hold:

$$P_L(X + \Delta)P_R^T = L \cdot R, \quad \|\Delta\| \leq \varepsilon_x \|X\|, \quad r_{ii} \neq 0, l_{ii} = 1 \text{ for every } i. \quad (\text{A1.1})$$

Let G be the *inverse* of X , *computed* by the B-method *via* GECP of X (that means: from the factors P_L, P_R, L, R). Then the matrix G satisfies the relations:

$$G + \Delta' = (X + \Delta)^{-1}, \quad \|\Delta'\| \leq \varepsilon_g \|G\|, \quad (\text{A1.2})$$

where $\varepsilon_x, \varepsilon_g$ are modest multiples of ν (the computing precision).

Proof. We will use here the ∞ -norm of matrices: $\|\cdot\| = \|\cdot\|_\infty$. Not lessening the generality of considerations let assume $P_L = I = P_R$. Let introduce the matrices

$$D \stackrel{\text{df}}{=} \text{diag}(r_{ii}), \quad U \stackrel{\text{df}}{=} D^{-1}R = [u_{ij}]. \quad (\text{A1.3})$$

The GECP process guarantees the relations for every i, j :

$$l_{ii} = u_{ii} = 1, \quad |l_{ij}| \leq 1, \quad |u_{ij}| \leq 1,$$

what implies the bounds:

$$\|L\| \leq n, \quad \|U\| \leq n, \quad \|L^{-1}\| \leq 2^{n-1}, \quad \|U^{-1}\| \leq 2^{n-1}. \quad (\text{A1.4})$$

Let present the B-method as following two assignment-statements:

$$V := R^{-1}, \quad G := V * L^{-1}. \quad (\text{A1.5})$$

Let $\mathbf{v}_i^T, \mathbf{g}_i^T$ be the i -th rows of V and G , respectively. Row-wise implementation of (A1.5) amounts to solving the following triangular equations:

$$\mathbf{v}_i^T R \stackrel{!}{=} \mathbf{e}_i^T, \quad \mathbf{g}_i^T L \stackrel{!}{=} \mathbf{v}_i^T \quad (i = 1, \dots, n),$$

where \mathbf{e}_i^T is the i -th row of the identity matrix. The computed solutions $\mathbf{v}_i, \mathbf{g}_i$ of these equations satisfy the equalities

$$\mathbf{v}_i^T (R + \delta R_i) = \mathbf{e}_i^T, \quad \mathbf{g}_i^T (L + \delta L_i) = \mathbf{v}_i^T, \quad (\text{A1.6})$$

where the *perturbation matrices* $\delta R_i, \delta L_i$ (equivalent to rounding-errors in the solving algorithms) are bounded:

$$|\delta R_i| \leq \nu c |R|, \quad |\delta L_i| \leq \nu c |L|. \quad (\text{A1.7})$$

($c \approx 1$, if “inner products” are cumulated on higher-precision variable, otherwise $c = n$.)

Let rewrite the equalities (A1.6) in the form

$$\mathbf{v}_i^T R = \mathbf{e}_i^T (I + \Phi_i)^{-1}, \quad \mathbf{g}_i^T (I + \Psi_i) = \mathbf{v}_i^T L^{-1}, \quad (\text{A1.8})$$

where, with $\delta U_i \stackrel{\text{df}}{=} D^{-1} \delta R_i$, see (A1.3),

$$\Phi_i \stackrel{\text{df}}{=} R^{-1} \delta R_i = U^{-1} \delta U_i, \quad \Psi_i \stackrel{\text{df}}{=} \delta L_i L^{-1}. \quad (\text{A1.9})$$

(We assume $\|\Phi_i\| < \frac{1}{2}$, since $\|\Phi_i\| \leq \nu c n 2^{n-1}$, see (A1.4), is for large n practically always a severe overbound.)

All row-equalities (A1.8) can be presented in the matrix form:

$$VR = I - \hat{\Phi}, \quad G + \Delta'_1 = VL^{-1}, \quad (\text{A1.10})$$

where [using the equality $(I + \Phi)^{-1} = I - \Phi(I + \Phi)^{-1}$] the i -th row of $\hat{\Phi}$ is equal to $\mathbf{e}_i^T \Phi_i (I + \Phi_i)^{-1}$, and the i -th row of Δ'_1 is equal to $\mathbf{g}_i^T \Psi_i$. From (A1.4), (A1.7), (A1.9) follow the bounds

$$\|\hat{\Phi}\| \leq \varepsilon_1 (1 - \varepsilon_1)^{-1}, \quad \|\Delta_1\| \leq \varepsilon_1 \|G\|, \quad \varepsilon_1 \stackrel{\text{df}}{=} \nu c n 2^{n-1}. \quad (\text{A1.11})$$

From (A1.1), (A1.10), (A1.11) we obtain ultimately

$$G + \Delta' = (X + \Delta)^{-1}, \quad \|\Delta'\| < \nu c n 2^n \|G\|, \quad (\text{A1.12})$$

where $\Delta' \stackrel{\text{df}}{=} \Delta'_1 + \hat{\Phi} (I - \hat{\Phi})^{-1} (G + \Delta'_1)$, what completes the proof. \square

Remarks A.1.

- (i) Relations (A1.1) are satisfied for any sufficiently well-conditioned matrix X . But simple modification of GECP guarantees (A1.1) (with ε_x being a modest multiple of ν) for any matrix $X \neq 0$. This allows us to apply the HS-process also for such matrices.
- (ii) Let's note that $\|\Delta'\| \approx \max_i \|\mathbf{g}_i^T \delta L_i L^{-1} + \mathbf{e}_i^T U^{-1} \delta U_i G\|$ and that the bounds (A1.4) on $\|L^{-1}\|$ and $\|U^{-1}\|$ are for larger n practically never approached. Hence in most cases (A1.12) is a severe *overbound* on $\|\Delta'\|$. We can expect that ε_g in (A1.2) is practically always a modest multiple of ν .
- (iii) In [10, pp. 110–111] Wilkinson proves the NC-property (A1.2) of the matrix G , computed *via* GEPP by the A-method, see [2], under assumptions that all involved triangular systems are *solved to high accuracy*. Since in the case of GECP this condition is always satisfied hence theorem A.1 is valid also for the A-method.