

The Reconstruction of Some 3D Convex Polyominoes from Orthogonal Projections

Maciej Gębala

Institute of Mathematics
Wrocław University of Technology
email: mgc@im.pwr.wroc.pl

Discrete Tomography

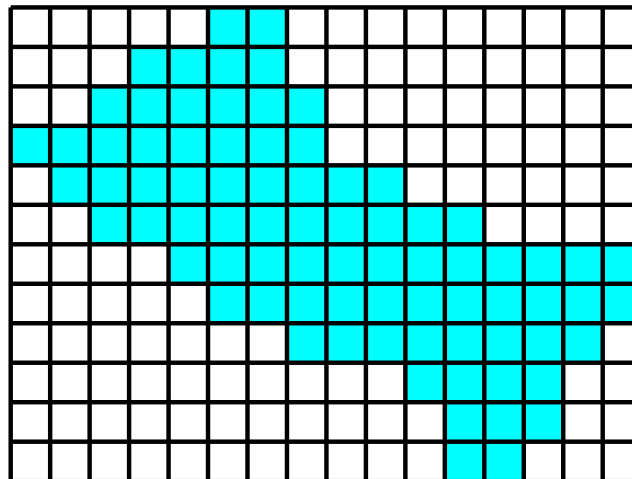
Basic Problem:

Reconstruction of finite point sets that are accessible only through some of their discrete X-rays.

2D Convex Polyomino

Set of cells S in a $m \times n$ grid with properties:

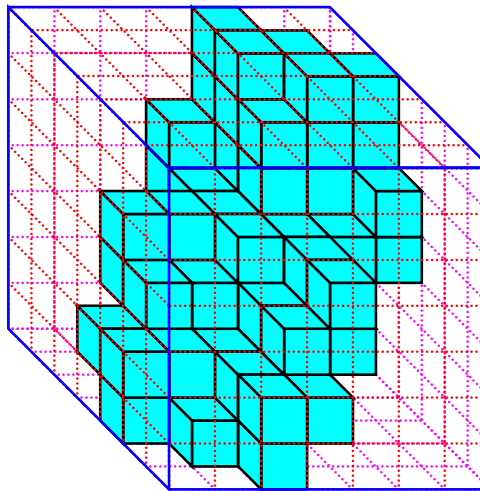
- S is connected,
- S in each row and each column is connected.



3D Convex Polyomino

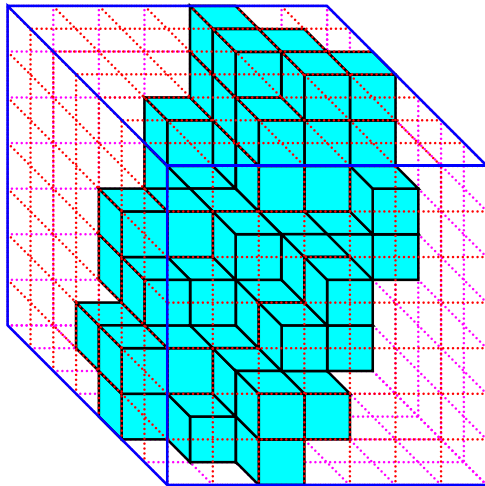
Set of cells S in $n \times n \times n$ grid with properties:

- S is connected,
- each two-dimensional orthogonal section of the grid contains 2D convex polyomino.



Orthogonal Projections

Orthogonal projections - the number of cells in each bar of the grid that belongs to S .



Top projections P_T

0	0	0	2	3	1	0
0	0	0	1	4	2	1
0	2	2	2	3	4	3
1	2	3	3	2	1	2
2	2	2	3	2	0	0
0	1	2	2	2	0	0
0	0	2	1	0	0	0

Front projections P_F

0	0	0	0	2	2	1
0	0	0	0	3	3	2
0	0	0	1	3	1	2
0	1	1	4	5	2	1
0	2	2	4	3	0	0
2	2	4	4	0	0	0
1	2	4	1	0	0	0

Side Projections P_S

1	2	2	0	0	0	0
2	3	3	0	0	0	0
2	1	3	1	0	0	0
1	2	5	4	1	1	0
0	0	3	4	2	2	0
0	0	0	4	4	2	2
0	0	0	1	4	2	1

Reconstruction Problem

Given three assigned matrices:

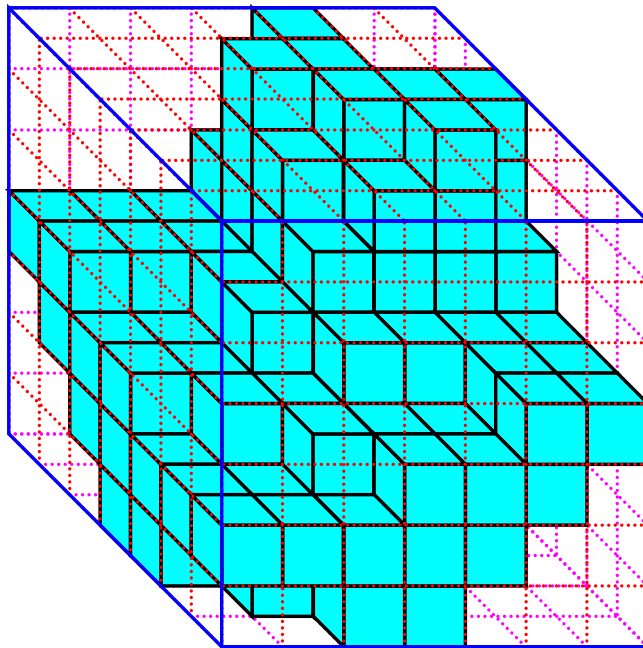
$$P_F, P_S, P_T \in \{0, \dots, n\}^{n \times n},$$

we examine whether there exists at least one convex polyomino S with adequate orthogonal projections.

Full 3D Convex Polyomino

3D Convex Polyomino with at least one matrix of projections does not contain zeros (*full matrix*).

We determine that P_T is a full matrix.



P_T

1	1	1	2	3	1	1
2	2	2	4	3	1	1
2	1	2	3	3	3	3
3	2	2	3	3	4	2
2	3	2	2	2	2	1
1	2	3	2	1	2	1
1	1	2	3	2	2	1

P_F

0	0	0	1	3	2	1
0	0	0	2	4	4	3
0	0	0	4	4	2	2
2	2	4	5	2	4	4
3	4	5	4	3	3	0
5	4	3	2	1	0	0
2	2	2	1	0	0	0

P_S

1	2	3	1	0	0	0
3	4	4	2	0	0	0
2	2	4	4	0	0	0
4	4	2	5	4	2	2
0	3	3	4	5	4	3
0	0	1	2	3	4	5
0	0	0	1	2	2	2

Our Result

Theorem. The problem of the reconstruction of full convex 3D polyominoes from orthogonal projections has the complexity $O(n^7 \log n)$.

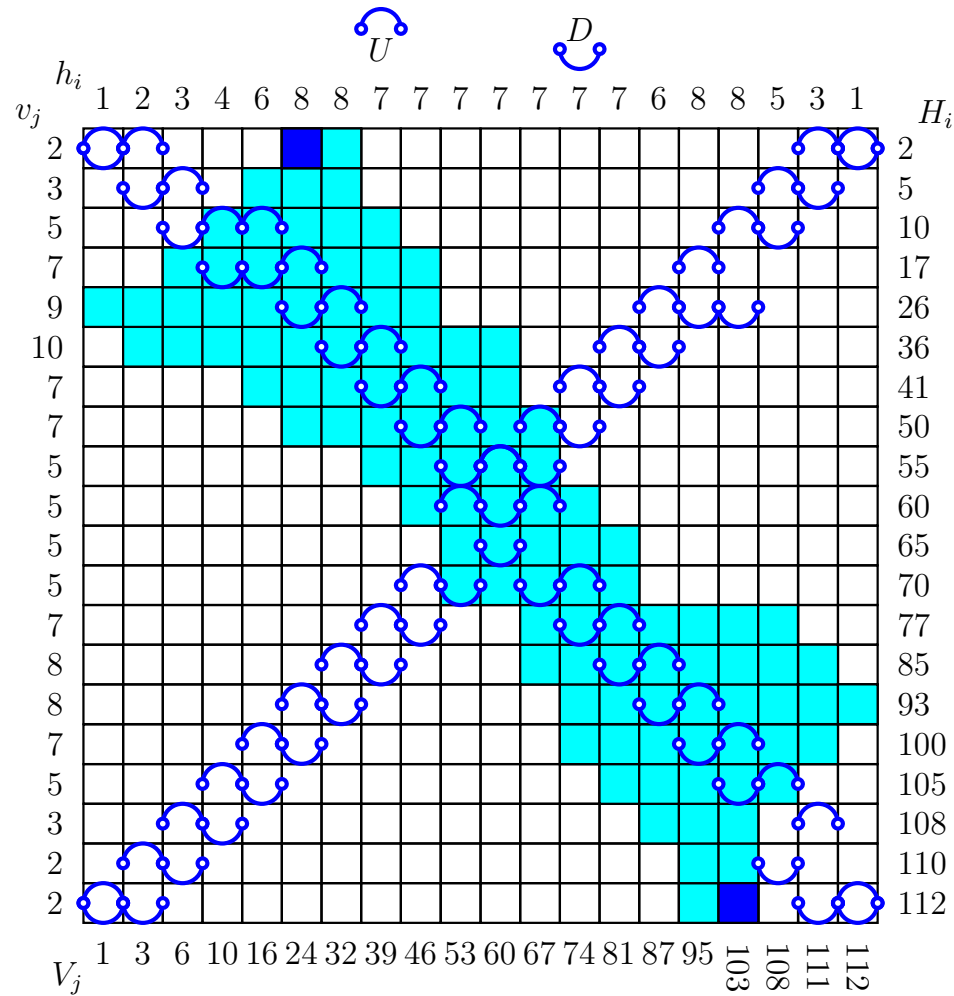
Properties of 2D Convex Polyom.

- auxiliary values $D_j^{\searrow}, U_j^{\searrow}, D_j^{\swarrow}, U_j^{\swarrow}$ (easy computable)
- $\forall_{j \in \{1, \dots, n\}} U_j^{\searrow} \leq D_j^{\searrow} \quad \wedge \quad U_j^{\swarrow} \leq D_j^{\swarrow}.$
- $\forall_{j \in \{1, \dots, n-1\}} D_j^{\searrow} + 1 \geq U_{j+1}^{\searrow} \quad \wedge \quad D_{j+1}^{\swarrow} + 1 \geq U_j^{\swarrow}.$

Lemma. Let cells $[1, p_1]$ and $[n, p_2]$ belong to the convex polyomino S . Then

1. if $p_1 \leq p_2$ then cells $[1, p_1], \dots, [D_{p_1}^{\searrow}, p_1]$, and $[U_j^{\searrow}, j], \dots, [D_j^{\searrow}, j] \quad \forall_{j \in \{p_1+1, \dots, p_2-1\}}$, and $[U_{p_2}^{\searrow}, p_2], \dots, [n, p_2]$ also belong to the S , or
2. if $p_1 \geq p_2$ then cells $[U_{p_2}^{\swarrow}, p_2], \dots, [n, p_2]$, and $[U_j^{\swarrow}, j], \dots, [D_j^{\swarrow}, j] \quad \forall_{j \in \{p_2+1, \dots, p_1-1\}}$, and $[D_{p_1}^{\swarrow}, p_1], \dots, [1, p_1]$ also belong to the S .

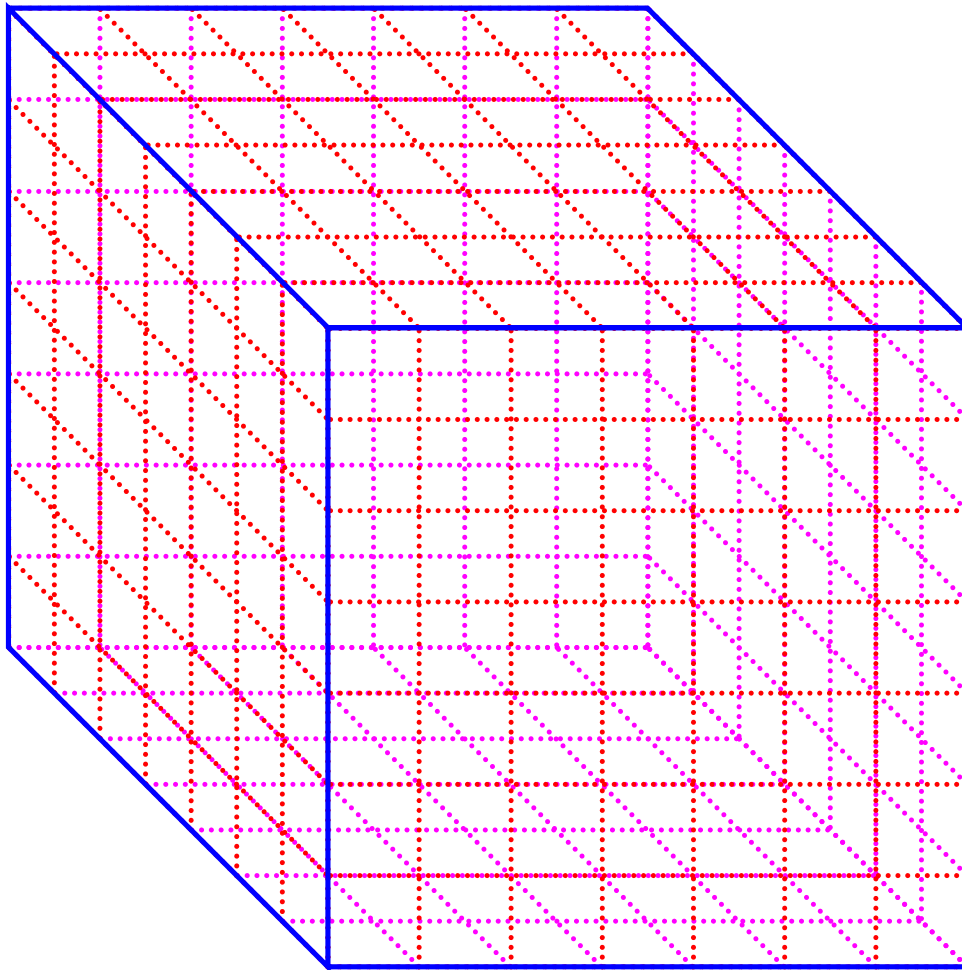
Example of Properties



Algorithm

1. Compute start positions
 - (a) select one cell in each corner vertical bar
 - (b) from Lemma compute positions of cells belonging to S in both lateral slices of the grid
 - (c) from Lemma compute positions of cells belonging to S in vertical slices orthogonal to slices from (b)
2. Perform Filling Procedure

Example of Reconstruction



P_T

1	1	1	2	3	1	1
2	2	2	4	3	1	1
2	1	2	3	3	3	3
3	2	2	3	3	4	2
2	3	2	2	2	2	1
1	2	3	2	1	2	1
1	1	2	3	2	2	1

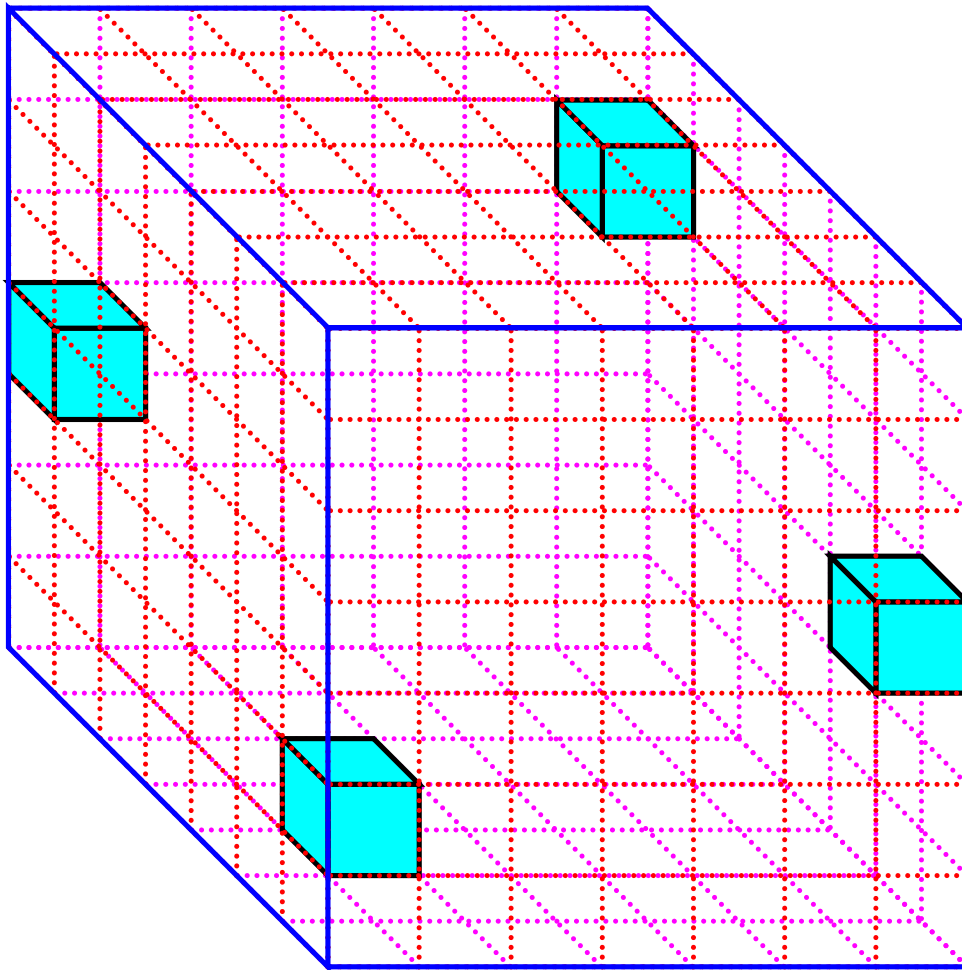
P_F

0	0	0	1	3	2	1
0	0	0	2	4	4	3
0	0	0	4	4	2	2
2	2	4	5	2	4	4
3	4	5	4	3	3	0
5	4	3	2	1	0	0
2	2	2	1	0	0	0

P_S

1	2	3	1	0	0	0
3	4	4	2	0	0	0
2	2	4	4	0	0	0
4	4	2	5	4	2	2
0	3	3	4	5	4	3
0	0	1	2	3	4	5
0	0	0	1	2	2	2

Example of Initial Positions



P_T

1	1	1	2	3	1	1
2	2	2	4	3	1	1
2	1	2	3	3	3	3
3	2	2	3	3	4	2
2	3	2	2	2	2	1
1	2	3	2	1	2	1
1	1	2	3	2	2	1

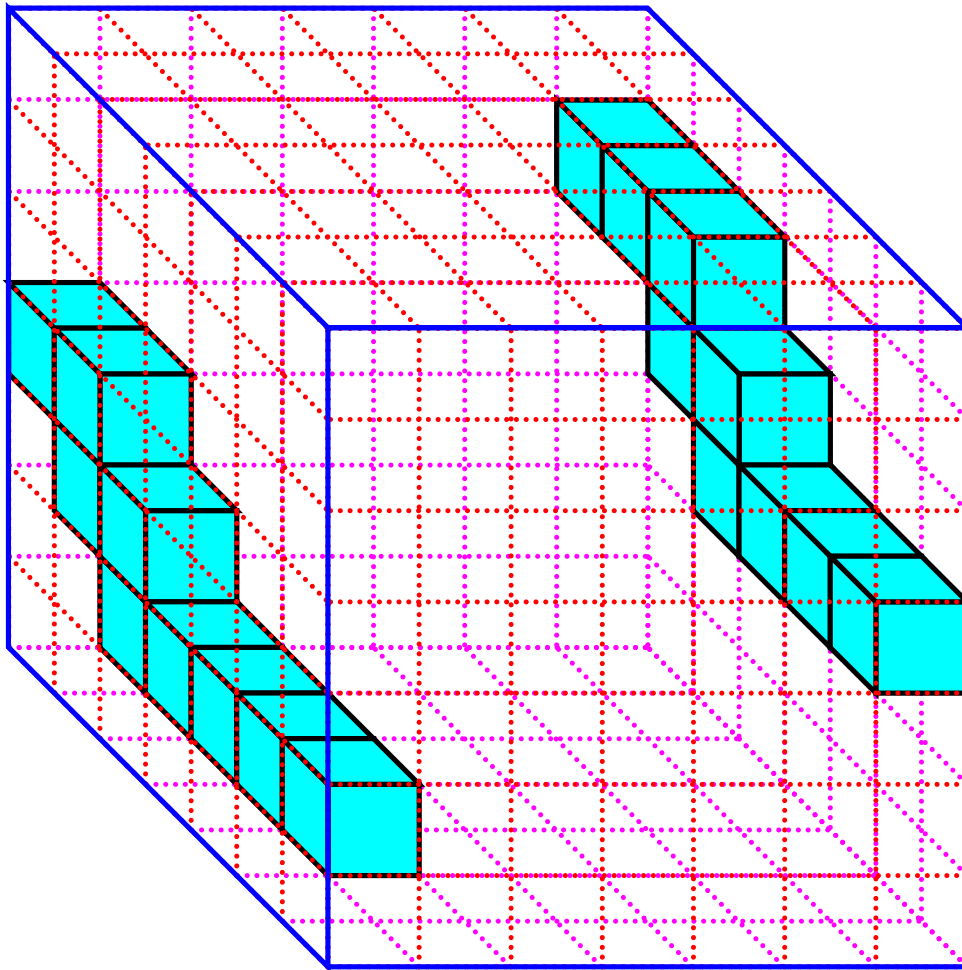
P_F

0	0	0	1	3	2	1
0	0	0	2	4	4	3
0	0	0	4	4	2	2
2	2	4	5	2	4	4
3	4	5	4	3	3	0
5	4	3	2	1	0	0
2	2	2	1	0	0	0

P_S

1	2	3	1	0	0	0
3	4	4	2	0	0	0
2	2	4	4	0	0	0
4	4	2	5	4	2	2
0	3	3	4	5	4	3
0	0	1	2	3	4	5
0	0	0	1	2	2	2

Example of Using the Lemma



P_T

1	1	1	2	3	1	1
2	2	2	4	3	1	1
2	1	2	3	3	3	3
3	2	2	3	3	4	2
2	3	2	2	2	2	1
1	2	3	2	1	2	1
1	1	2	3	2	2	1

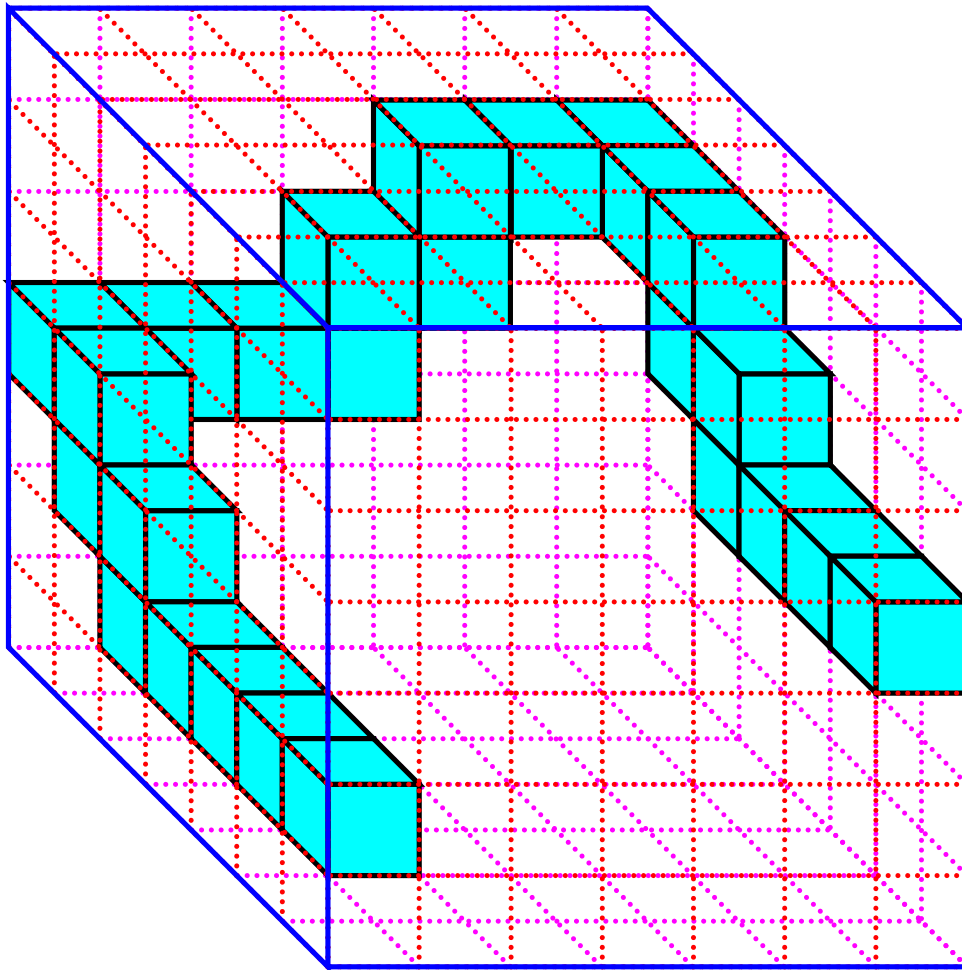
P_F

0	0	0	1	3	2	1
0	0	0	2	4	4	3
0	0	0	4	4	2	2
2	2	4	5	2	4	4
3	4	5	4	3	3	0
5	4	3	2	1	0	0
2	2	2	1	0	0	0

P_S

1	2	3	1	0	0	0
3	4	4	2	0	0	0
2	2	4	4	0	0	0
4	4	2	5	4	2	2
0	3	3	4	5	4	3
0	0	1	2	3	4	5
0	0	0	1	2	2	2

Example of Using the Lemma



P_T

1	1	1	2	3	1	1
2	2	2	4	3	1	1
2	1	2	3	3	3	3
3	2	2	3	3	4	2
2	3	2	2	2	2	1
1	2	3	2	1	2	1
1	1	2	3	2	2	1

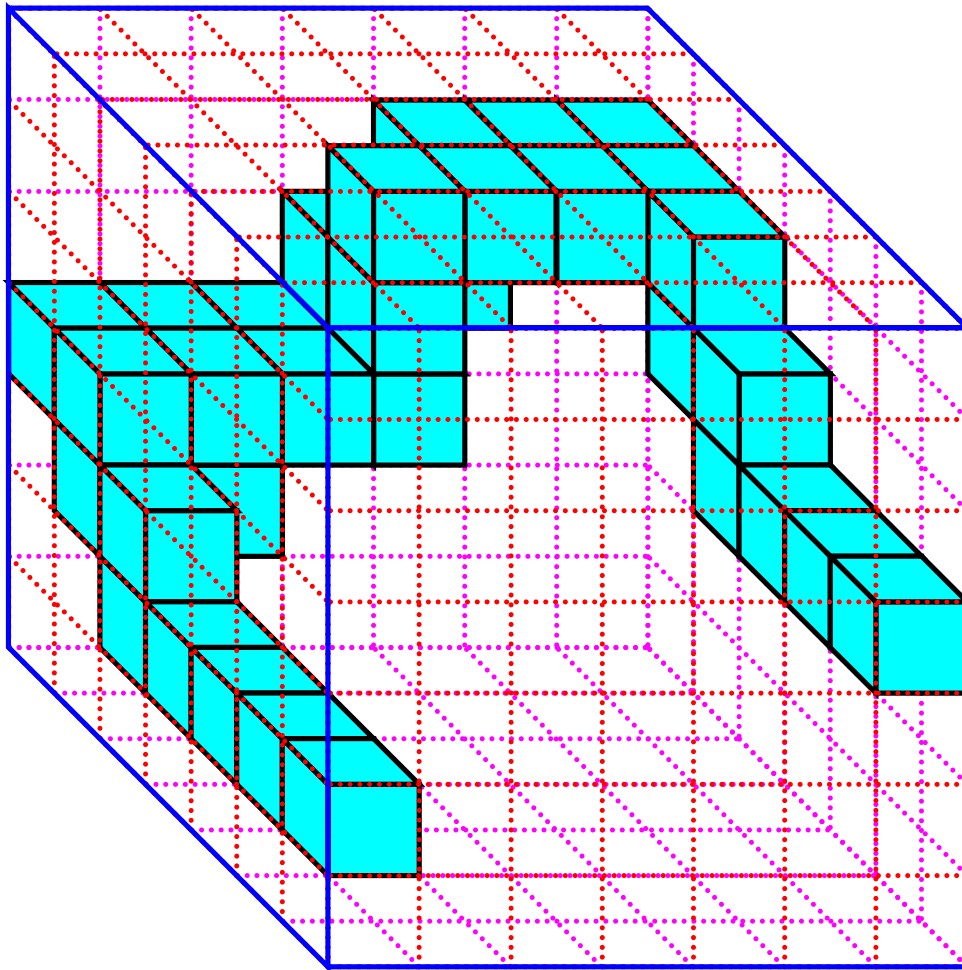
P_F

0	0	0	1	3	2	1
0	0	0	2	4	4	3
0	0	0	4	4	2	2
2	2	4	5	2	4	4
3	4	5	4	3	3	0
5	4	3	2	1	0	0
2	2	2	1	0	0	0

P_S

1	2	3	1	0	0	0
3	4	4	2	0	0	0
2	2	4	4	0	0	0
4	4	2	5	4	2	2
0	3	3	4	5	4	3
0	0	1	2	3	4	5
0	0	0	1	2	2	2

Example of Using the Lemma



P_T

1	1	1	2	3	1	1
2	2	2	4	3	1	1
2	1	2	3	3	3	3
3	2	2	3	3	4	2
2	3	2	2	2	2	1
1	2	3	2	1	2	1
1	1	2	3	2	2	1

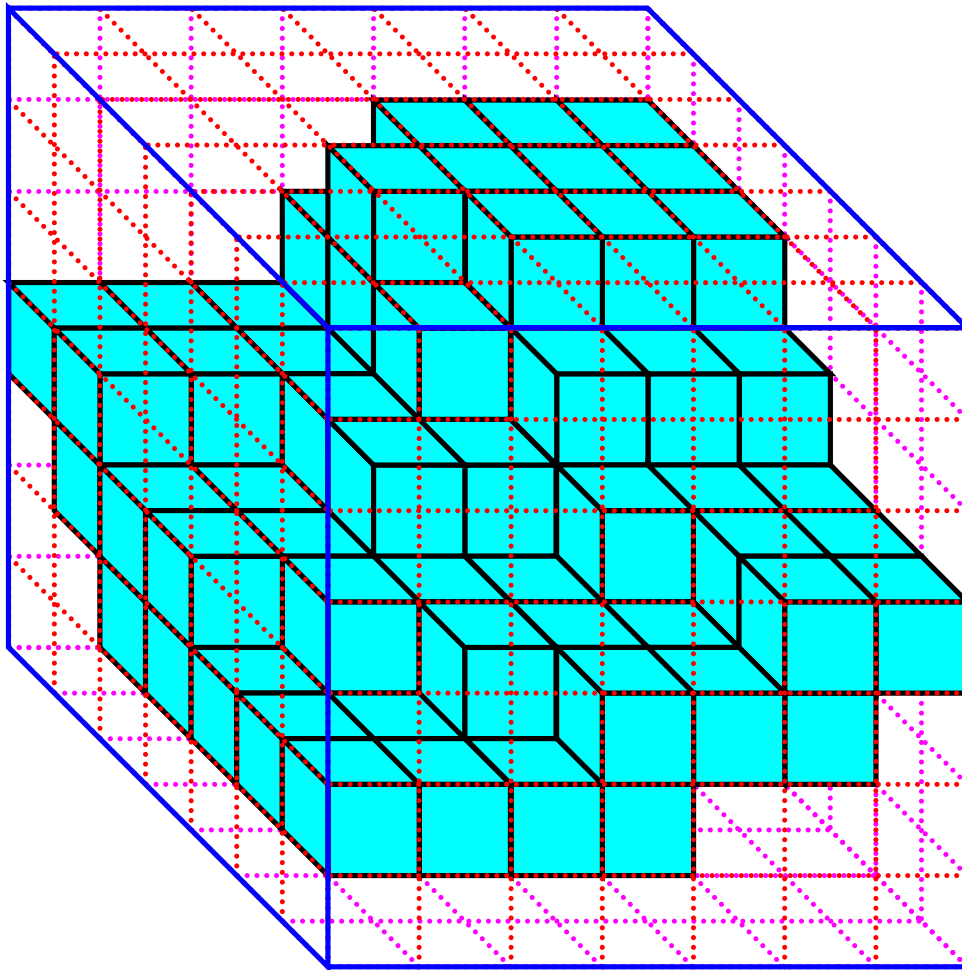
P_F

0	0	0	1	3	2	1
0	0	0	2	4	4	3
0	0	0	4	4	2	2
2	2	4	5	2	4	4
3	4	5	4	3	3	0
5	4	3	2	1	0	0
2	2	2	1	0	0	0

P_S

1	2	3	1	0	0	0
3	4	4	2	0	0	0
2	2	4	4	0	0	0
4	4	2	5	4	2	2
0	3	3	4	5	4	3
0	0	1	2	3	4	5
0	0	0	1	2	2	2

Example of Start Positions



P_T

1	1	1	2	3	1	1
2	2	2	4	3	1	1
2	1	2	3	3	3	3
3	2	2	3	3	4	2
2	3	2	2	2	2	1
1	2	3	2	1	2	1
1	1	2	3	2	2	1

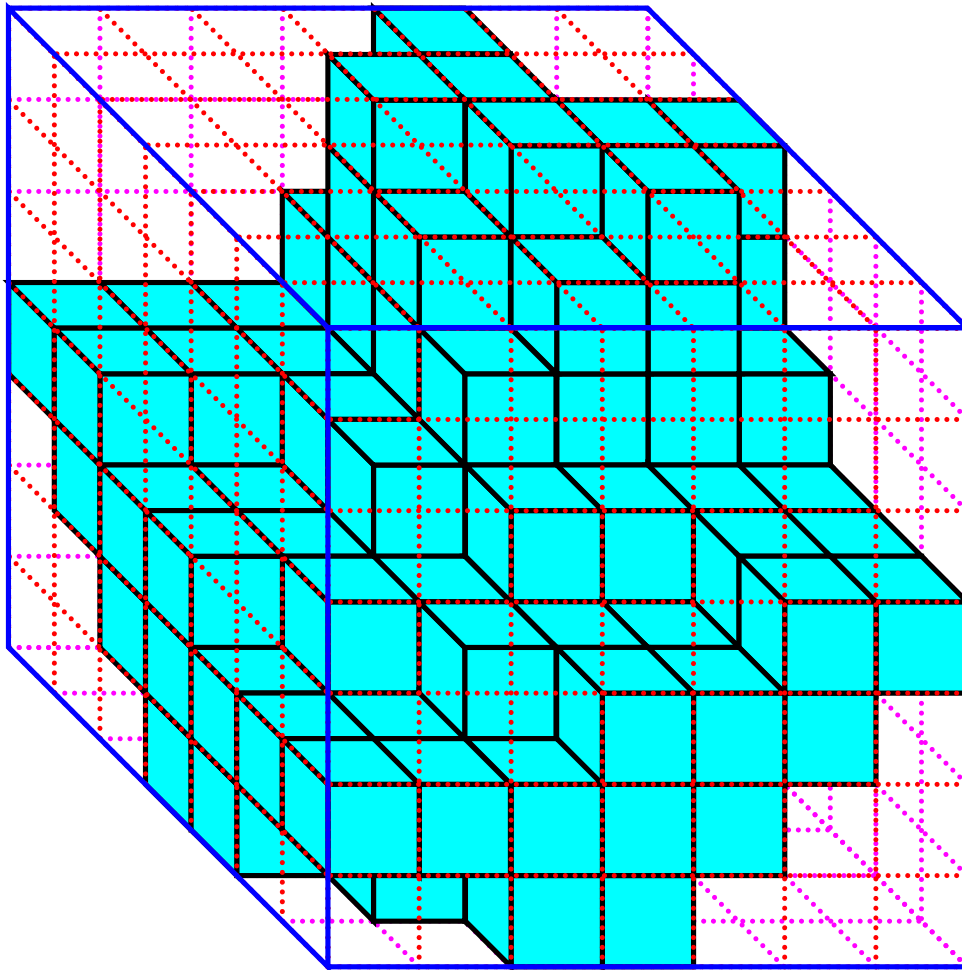
P_F

0	0	0	1	3	2	1
0	0	0	2	4	4	3
0	0	0	4	4	2	2
2	2	4	5	2	4	4
3	4	5	4	3	3	0
5	4	3	2	1	0	0
2	2	2	1	0	0	0

P_S

1	2	3	1	0	0	0
3	4	4	2	0	0	0
2	2	4	4	0	0	0
4	4	2	5	4	2	2
0	3	3	4	5	4	3
0	0	1	2	3	4	5
0	0	0	1	2	2	2

Example of 3D Convex Polyomino



P_T

1	1	1	2	3	1	1
2	2	2	4	3	1	1
2	1	2	3	3	3	3
3	2	2	3	3	4	2
2	3	2	2	2	2	1
1	2	3	2	1	2	1
1	1	2	3	2	2	1

P_F

0	0	0	1	3	2	1
0	0	0	2	4	4	3
0	0	0	4	4	2	2
2	2	4	5	2	4	4
3	4	5	4	3	3	0
5	4	3	2	1	0	0
2	2	2	1	0	0	0

P_S

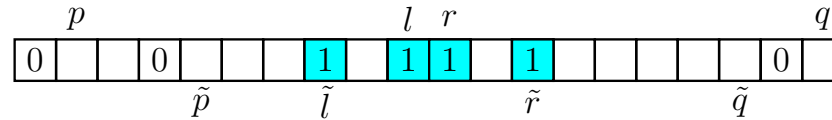
1	2	3	1	0	0	0
3	4	4	2	0	0	0
2	2	4	4	0	0	0
4	4	2	5	4	2	2
0	3	3	4	5	4	3
0	0	1	2	3	4	5
0	0	0	1	2	2	2

Filling Procedure - properties

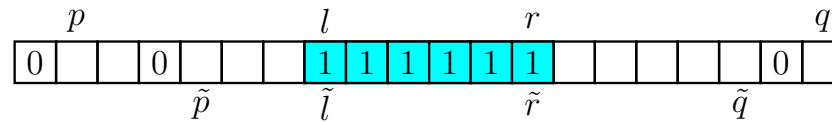
- It returns *fail* if 3D convex polyomino with computing start positions and adequate matrices of projections does not exist.
- Otherwise it returns **3D convex polyomino**.
- It works correctly if start positions contain at least one cell of polyomino in each vertical bar.

Filling Procedure - the main idea

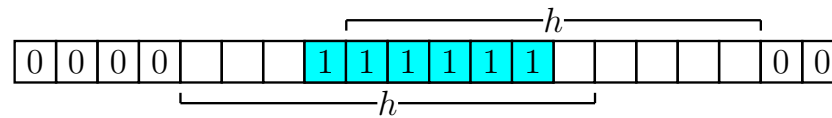
$h = 10$ - the value of the projection for the bar



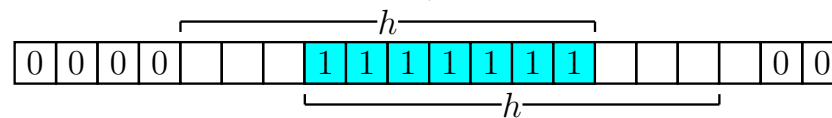
operation \oplus - the integration of the block of 1's



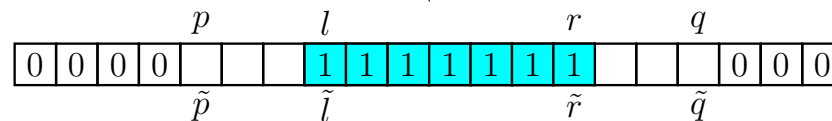
operation \ominus - the integration of the two final blocks of 0's



operation \otimes - the expansion of the block of 1's



operation \odot - the expansion of the two final blocks of 0's



Complexity

- n^4 – the number of different initial positions;
- $O(n^2)$ – the cost of computing of the start positions;
- $O(n^3 \log n)$ – the cost of the Filling Procedure.

Complexity of the Algorithm – $O(n^4 \cdot n^3 \log n)$.