Self-stabilization

Algo 21



Types of errors:

--- failures of comm

-- Byzantine communication





Self stabilization

End of problems <a>[? transition to a proper shape

Token ring communication





Algorithm 13.3 Self-stabilizing Token Ring
1: if
$$v = v_0$$
 then
2: if $S(v) = S(p)$ then
3: $S(v) := S(v) + 1 \pmod{n}$
4: end if
5: else
6: $S(v) := S(p)$
7: end if









Algorithm 13.5 Self-stabilizing MIS

Require: Node IDs

Every node v executes the following code:

- 1: do atomically
- 2: Leave MIS if a neighbor with a larger ID is in the MIS
- 3: Join MIS if no neighbor with larger ID joins MIS
- 4: Send (node ID, MIS or not MIS) to all neighbors

5: end do



Transformation of deterministic algorithms into selfstabilization



messages from radius d! run algorithm - emulate

Randomized case



Advance example of self-stabilization

Every evening:

- -- a voter calls the friends
- -- the friends give their recommendations
- -- the voter changes preference according to majority of recommendations



- Is eventually everybody voting for the same party?
- Will each citizen eventually stay with the same party?
- Will citizens that stayed with the same party for some time, stay with that party forever?
- And if their friends also constantly root for the same party? No.



Theorem 13.7 (Dems & Reps). Eventually every citizen is rooting for the same party every other day.

Day t: supporter of Dems

Bad out-edges Opem Good out-edges De Jen Rep Dem Der Jew JC~ 2

Day t:

Day t+1:

b bad out-edges from step t

g good out-edges from step t

recommendations for Reps

7 recommendations for Dems



The voter roots Dems again

bad in-edges at step t+1 = # bad out-edges at step t



Case: g<b

The voter roots REPs

bad in-edges at step t+1 = g
bad out-edges at step t = b



Total number of bad edges in the graph

Total number of bad edges in the graph





/em ъ 0 b 6 6 Ю-Ť ð





