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# Security and Cryptography 2021 Mirosław Kutyłowski

#### grading criteria:

- up to 50 points from lecture (exam), up to 50 points from dr Kubiak (project...)
- the lecture at least 30% of 50 points must be earned to pass
- sum of points  $\Rightarrow$  the final grade, 3.0:  $\geq$ 40 points , 5.0  $\geq$  80 points
- exam requires problem solving, memorizing facts is unnecessary

skills to be learned: developing end-to-end security systems, flawless in the real sense!

- presence: obligatory during the lectures
- exam date and form: subject to the situation
- place: 11:15-13 Wednesday, 11:15-13 Friday, MS Teams

adjustments possible in order to ease logistics problems

#### grading system used last year:

- for each "chapter" consisting of a specific topic some verification of skills of the students
- possible verification forms:
  - i an assignment by myself (some concrete task/problem to be solved at home and returned within e.g. 1 week)
  - ii written exam with ePortal

answer a problem to be typed in or solved on paper, jpeg to be uploaded within, say, 15 minutes

• IPR taken very seriously

#### **Online materials**:

• available on my webpage

https://kutylowski.im.pwr.wroc.pl/lehre/cs21/

- ePortal will be used for 1-1 communiaction with the students (as it keeps a history of each conversation),
- tests no decision about the platform yet. Subject to: stability and reliability of these tools

#### Contact:

• during the lecture: mute yourself, when having a question please unmute and switch on your video (the group is small enough to do it)

WCS(

- email: yes, but assignments etc over ePortal
- the phone at Politechnika no!
- MS Teams for conf calls
- Signal as second independent channel?

# I. FAILURE EXAMPLES TO LEARN FROM I.1. PKI for Signing Digital Documents

#### **PKI - Public Key Infrastructure**

- strong authentication of digital documents with digital signatures seems to be possible
- in fact we get an evidence that the holder of a private key has created a signature
- who holds the key? PKI has to provide a certified answer to this question
- PKI is not a cryptographic solution it is an organizational framework (using some crypto key <>>> physical person [Doardin P] zanfany podpis osobisty tools)

PESEL

> root CA

#### PKI, X.509 standard

- a certificate binds a public key with an ID of its alleged owner,
- a couple of other fields, like validity date, key usage, certification policy, ...
- certificate signed by CA (Certification Authority)
- tree of CA's (or a directed acyclic graph), with roots as "roots of trust"
- status of a certificate may change revocation
- checking status methods: CRL, OCSP

#### reasons for PKI failure:

- a nice concept of digital signatures but
  - 1. big infrastructure required:
    - -. substantial cost and effort
    - -. long time planning needed (so possible in China, but not in Europe)
    - -. unclear financial return
- 2. scope of necessary coordination,
  - -. in order to work must be designed at least for the Common Market
  - -. example of killing the concept: link to certification policy in Polish

- 3. lack of interoperability (sometimes as business goal)
  - -. companies make efforts to eliminate competition
  - -. standarization may be focused on securing market shares
  - -. a long process  $\ldots$
- 4. necessary trust in roots
  - –. how do you know that the root is honest?
- 5 registration: single point of fraud, (e.g. with fake breeding documents)
  - -. once you get a certificate you may forge signatures
- 6. responsibility of CA 5. responsibility of CA 5. responsibility 6. responsibility 200 PL 7. cost who will pay? For the ord user the initial cost is too high
  - 7. cost who will pay? For the end user the initial cost is too high.
    - -. certificates are too expensive for just a few signatures (at least initially)
  - 8. legal strength of signatures
    - -. if scheme broken or signing devices turn out to be insecure you are anyway responsible for the signatures. After revocation only the new signatures invalid

( podpis osobist, on

SIGNO

9. unsolved problem of revocation: possible to check the status in the past but not now

reason: mismatch of requirements and interests with the designed solution

".(.but there nothing one can do about it." )- this is false

- Smart-ID project, Estonia (clever RSA-like solution, mediated signatures, no CRL, OCSP needed)
- SPKI idea (source centric certification), suicide notes, certificates of health



#### before Smart-ID in Estonia

- personal ID smart cards, implements RSA signature of the owner
- certificate of BSI for Infineon chip and software
- Czech colleagues from Brno found that the RSA keys generated so that the old attacks work

not junne against C.A.

- an implementation bug or a trapdoor
- all smart cards had to be updated

In

#### Smart-ID

1. RSA:

-. "RSA" where n is a product of two RSA numbers

- -. the same algebra no difference seen unless you factorize n
- -. but secret keys distributed between the card and a mediator server
- $-.\ nobody \ has \ full \ knowledge \ of \ the \ secret \ keys$
- 2. links between consecutive signatures (to be checked by the mediator server)

3. revocation by blacklisting on the server





# I.2. Clickjacking on Android

#### **Overlay mechanism:**

- apps are separated in their sandboxes security design mechanism
- all apps display informations on the screen at the same time they are overlays
- overlays:
  - require Android permissions
  - clickable or paththrough
  - opaque or transparent
  - combining: parameter lpha defines weights: for each color of RGB the new pixel value is

$$\mathrm{old} \cdot \alpha + \mathrm{new} \cdot (1 - \alpha)$$



- basic clickjacking:
  - on top an opaque overlay with something innocent (game...)
  - button is in fact a batton but the overlay is paththrough at this place
  - below is an unvisible button of an attacked app

defense: Google's "obscure flag" - an app checks if at the moment of clicking there is an overlay above it



• context-hiding clickjacking: overlay covers everything but not the button

defense: Google's "hide overlays" (applicable only in case of settings etc as the users like overlays)



- examples of attacks:
  - $\rightarrow$  Google play: after installing the app it asks for ''open app'', but acceptance is for installing and opening something else
  - $\rightarrow$  Browser: cover the context and make the user click (e-voting?)
  - $\rightarrow$  gmail: prepare a message, cover it with overlay and ask for accepting "send" button
  - $\rightarrow$  whatsapp, we chat send messages, send SMS to chosen destinations (and learn attacked SIM subscriber number)
    - → Google Authenticator: "long click" copies a token to clipboard (and makes it available to other apps)
  - $\rightarrow$  Facebook, Tweeter: unprotected, possibilities to insert likes, send tweets, ...
  - $\rightarrow$  Lookout Mobile Security: 3 clicks and anti-virus protection disabled

- remedy?
  - no effective protection mechanism known
  - - architecture separates apps so it is impossible to ask what the other apps are showing
  - some overlays must be tolerated due to expectations of the users
  - pressure to run security critical apps (banking...)

#### **Clickshield:**

- requires minor changes in the Android framework
- concentrates on the central region (as on margin no critical buttons observed)
- attempts to check whether between the pixel of a target app and the screen final render pixel there is a nonmalicious relationship
- $\alpha$  render: fr = round( $\alpha \cdot ov + (1 \alpha) \cdot ta$ ) where ta=target app pixel value, ov=overlay..., fr=final render pixel value

login: student passvod: hics21

- $-\,$  given fr and ta we do not know  $\alpha$  and ov, but
  - choose two points and assume that the overlay is uniform. Then solve for  $\alpha$ :

$$\int \mathbf{fr}_1 = \alpha \mathbf{ov} + (1 - \alpha) \cdot \mathbf{ta}_1$$
$$\mathbf{fr}_2 = \alpha \cdot \mathbf{ov} + (1 - \alpha) \cdot \mathbf{ta}_2$$

- use this fixed  $\alpha$  to estimate ov over the whole screen. the outcome should be uniform in case of uniform verlays or those informing on margin (notifications)

Otherwise expect a context switching.

## I.3. Easy Fishing on Android

- mobile password managers:
  - associate app (package name) with a domain name
  - $-\,$  for a domain name associated with the app insert the user's credentials when the app accesses this URL
  - credentials are related to the domain name (facebook.com, etc) and not to the app
  - in fact: improves protection against fishing (difference between facebook and faceb00k detected, a human may make a mistake)

- Instant Apps: instead of downloading the whole app fetch only a small app that emulates the full version with accessing an URL
  - gets a full control over the screen e.g. it may hide browser's security information
- limitations: at most one app with the same package name on an Android device, and at most one in Play Store

- attack:
  - create an app
  - choose the package name so that the mapping points to the attacked domain
  - include developer's URL for Instant App purposes (it is not checked by the Play Store)
  - lure the user to run Instant App
    - $\rightarrow~$  the credentials will be included by the Password Manager
    - $\rightarrow~$  the information will go to the developer's URL
    - → Instant App will show something different on the screen (information from password manager need not to be visible, browser information will be covered as well)
    - $\rightarrow$  the user should dislike the app and consequently the app will not be downloaded (no traces of forgery)

- mappings:
  - most important associations on a kind of whitelist
  - one-to-one would be more secure but frequently many apps to one domain
  - (insecure) heuristics:
    - 1 Keeper: finds a corresponding entry on play.google.com and takes "app developer website field", it autosuggests this name to the user, attacks: write malicious app "developer website field"

2 Dashlane: hardcopied 81 mappings, rest: autosuggestion heuristic based on at least
3 matching characters: xxx.face.yyy will be mapped to facebook.com

attack: use similar package names

- 3 Lastpass: translates directly (aaa.bbb.ccc to bbb.aaa), if not existing then consult from crowdsourced mapping (distributed database created by users where it is easy to inject something)
- 4 1Password: does not provide mapping but presents suggestions and enables the user
  a human to choose (and confuse FACEB00K with FACEBOOK)
- 5 Google Smart Lock burden of mapping to the developer (at the time not automated process) based on Digital Asset Links (on the website a list of permitted apps, authenticated with hashes of the signing key)

### I.4. Buying a system

**Problem:** somebody has to deploy a secure IT system, how to purchase it?

- problematic requirements according to BSI guide:
  - i **incomplete** forgetting some threats is common
  - ii not embedded: not corresponding really to the environment where the product has to be deployed
  - iii **implicit:** customer has in mind but the developer might be unaware of them
  - iv not testable: ambiguous, source of legal disputes, ...
  - v too detailed: unnecessary details make it harder to adjust the design

vi unspecified meaning: e.g. "protect privacy"

vii inconsistent: e.g. ignoring trade-offs

- specification-based purchasing process versus selection-based purchasing process
- the user is not capable of determining the properties of the product himself: too complicated, too specialized knowledge required, a single error makes the product useless
- specifications of concrete products might be useless for the customers hard to understand and compare the products
- informal specifications and descriptions, no access to crucial data



# I.5. Blind Trust

#### Idea - our dream situation:



Catacrypt don't wait for quantum computer, catastrophy is already there due to other reason