

**Probability and statistics, 2021, Computer Science
Algorithmics,
Undergraduate Course, Part II, lecturer: Mirosław
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**I. Generating Random Numbers for
a given probability distribution**

Chapter 5.2 in Byron

Goal:

**- simulations (e.g., pharma industry,
weather forecast, system testing ...)**

Weather simulations:

Simulations for new chemical products, pharmaceuticals:

Physical sources: examples:

1) Electronics (bistable)

3) quantum generators

3) noise

Problems of physical sources:

1) bias

2) memory: dependence on history

3) external influence

deterministic random number generators:

DRNG: NIST, recommendations,

architecture: PRNG(seed) yields: bits

DRNG:

basic property: not distinguishable from coin flipping

what does it mean?:

NIST tests

left-or-right game

secure PRNG:

**unpredictability
forwards:**

backwards:

realizations:

families of PRNG (based on residual arithmetic and algebraic expressions

$$a \cdot x^2 + b \cdot x + c \text{ mod } p$$

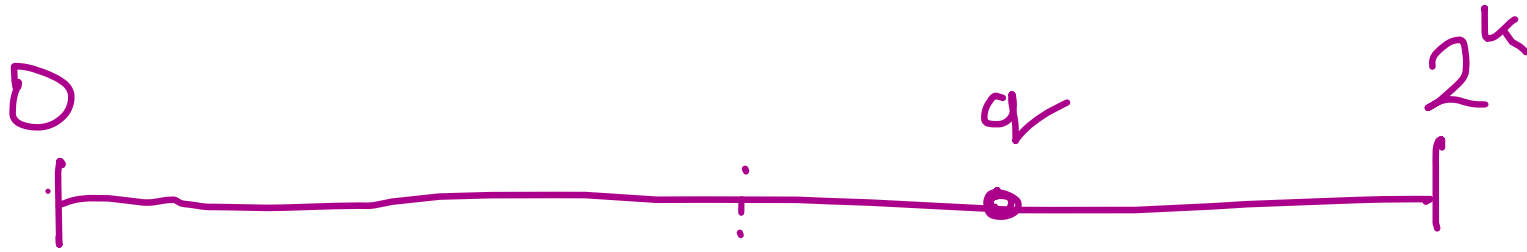
cryptographic generators: e.g. based on encryption

$$\text{trunc}(32, \text{Enc}_K(1)), \text{trunc}(32, \text{Enc}_K(2)), \dots$$

Problem: domain

We have a good generator for uniform distribution over n bit numbers

How to get a uniform distribution over integers in the range $[0, q)$?



1. choose $x \in [0, 2^k)$ at random
2. if $x \geq q$ then goto 1
3. output(x)

Problem: uniform versus non-uniform distribution

**all good PRNG resources deliver the output that is uniform over some interval
e.g.: 32 bit nonnegative integers**

Needed:

e.g. geometric distribution, Poisson, ...

single Bernoulli trial:

procedure:

1. choose u uniformly at random in $[0,1]$
2. if $u < p$ then output 0 else output 1

n Bernoulli trials, number of successes:

n times:

0: with probability p

1: with probability $1-p$

count the number of successes

stupid solution:

compute pbb according to formulas

...

n Bernoulli trials, number of successes:

n times:

0: with probability p

1: with probability $1-p$

count the number of successes

procedure (in MATLAB):

$n=20$; $p=0.34$;

$U=\text{rand}(n,1)$;

$X=\text{sum}(U<p)$

geometric distribution:

Bernoulli trials with pbb p of 0

output: the number of trials until 1 chosen

naive way: take mathematical formulas and then choose according to the probabilities

procedure (in MATLAB):

X=1;

while rand<p

X=X+1;

end;

X

arbitrary discrete distribution:

assume: n possible values, $p(i)$ -probability of the i th value

approach: for each $i \leq n$ compute

$$a_i = \sum_{j < i} p_j$$

the results saved in a data structure D

procedure:

1. $u = \text{random};$
2. with D find i such that $a_i \leq u < a_{i+1}$

Poisson distribution:

$$Pr(X=i) = e^{-\lambda} \cdot \frac{\lambda^i}{i!}$$

Mat

```
lambda = 5; % Parameter
U = rand; % Generated Uniform variable
i = 0; % Initial value
F = exp(-lambda); % Initial value, F(0)
while (U >= F); % The loop ends when U < F(i)
    F = F + exp(-lambda) * lambda^i/gamma(i+1);
    i = i + 1;
end;
X=i
```

the case of invertible CDF

Theorem 2 *Let X be a continuous random variable with cdf $F_X(x)$. Define a random variable $U = F_X(X)$. The distribution of U is $\text{Uniform}(0,1)$.*

PROOF: First, we notice that $0 \leq F(x) \leq 1$ for all x , therefore, values of U lie in $[0, 1]$. Second, for any $u \in [0, 1]$, find the cdf of U ,

$$\begin{aligned} F_U(u) &= P\{U \leq u\} \\ &= P\{F_X(X) \leq u\} \\ &= P\{X \leq F_X^{-1}(u)\} && \text{(solve the inequality for } X\text{)} \\ &= F_X(F_X^{-1}(u)) && \text{(by definition of cdf)} \\ &= u && (F_X \text{ and } F_X^{-1} \text{ cancel)} \end{aligned}$$

the case of invertible CDF - continuous distribution X

Procedure:

1. choose u uniformly at random in $[0,1]$

2. take

$$x := F_X^{-1}(u)$$

Example

Exponential distribution $F(x) = 1 - e^{-\lambda x}$

Procedure:

1. choose u uniformly at random in $[0,1]$

2. solve $u = 1 - e^{-\lambda x}$

that is $1 - u = e^{-\lambda x}$

$$\ln(1-u) = -\lambda x$$

$$x = -\ln(1-u)/\lambda$$

Or simply $x = -\ln(u)/\lambda$

Example –warning

Gamma distribution has complicated density function

$$F(t) = \int_0^t f(x)dx = \frac{\lambda^\alpha}{\Gamma(\alpha)} \int_0^t x^{\alpha-1} e^{-\lambda x} dx.$$

Inverting F ?

Workaround:

a random variable with Gamma distribution α is a sum of α independent random variables with exponential distribution

the case of invertible CDF - discrete distribution X

Procedure:

1. choose u uniformly at random in $[0,1]$

2. take

$$X := \min \{x : F(x) > u\}$$
$$\text{so: } X = F^{-1}(u)$$

Example: geometric distribution

$$F(x) = 1 - (1 - p)^x,$$

Procedure:

Find the smallest x such that

$$1 - (1 - p)^x > U$$

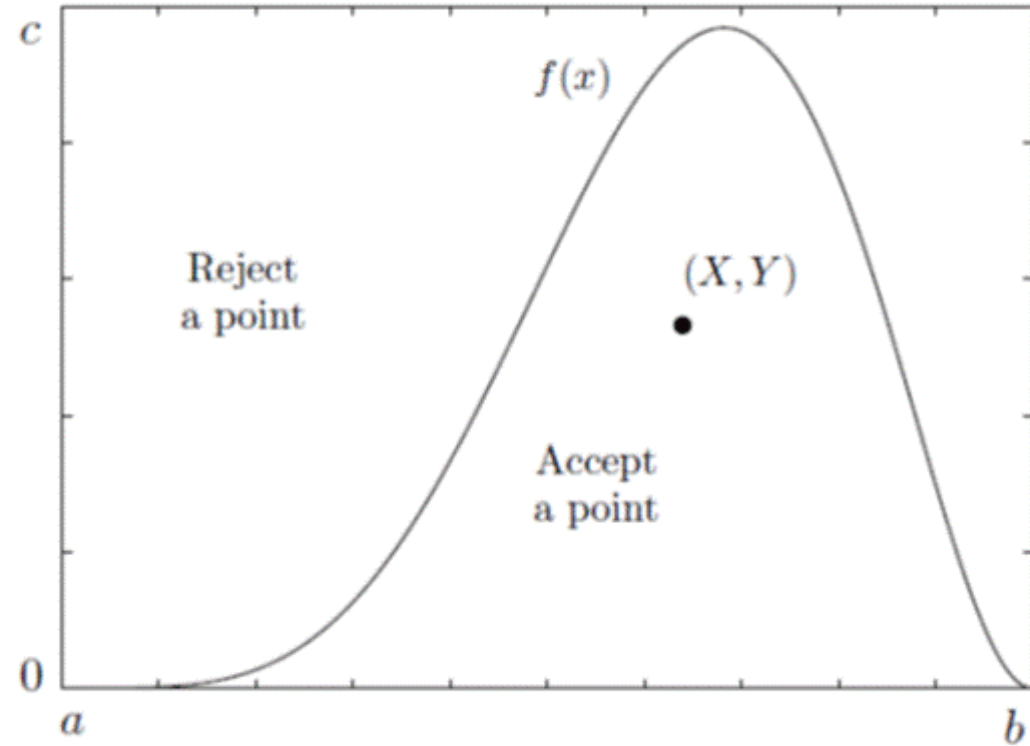
Solution:

$$X = \left\lceil \frac{\ln(1 - U)}{\ln(1 - p)} \right\rceil.$$

Rejection method:

1. sample X and Y uniformly at random
2. if $Y > f(X)$, then goto 1
3. Output X

then density of $X = f(X)$!



Pict. from Byron

Application:

distributions where density is computable (e.g. Beta distribution)

$$f(x) = \frac{\Gamma(\alpha + \beta)}{\Gamma(\alpha)\Gamma(\beta)} x^{\alpha-1} (1-x)^{\beta-1} \text{ for } 0 \leq x \leq 1.$$

but computing cdf is hard (numeric computations of the integral)

```
alpha=5.5; beta=3.1; a=0; b=1; c=2.5;
X=0; Y=c; % Initial values
while Y > gamma(alpha+beta)/gamma(alpha)/gamma(beta)...
    * X.^(alpha-1) .* (1-X).^(beta-1);
    U=rand; V=rand; X=a+(b-a)*U; Y=c*V;
end; X
```

Poisson distribution

An example of a clever approach tailored to the particular case

$$P(x) = e^{-\lambda} \frac{\lambda^x}{x!}, \quad x = 0, 1, 2, \dots$$

but.. also can be understood as the number of rare events in an interval of time, where the time between events is exponential

Pragmatic computation:

1. Obtain Uniform variables U_1, U_2, \dots from a random number generator.
2. Compute Exponential variables $T_i = -\frac{1}{\lambda} \ln(U_i)$.
3. Let $X = \max \{k : T_1 + \dots + T_k \leq 1\}$.