Probability and statistics, 2022, Computer Science Algorithmics, Undergraduate Course, Part II, lecturer: Mirosław Kutyłowski

3- Stochastic Processes

Stochastic process

- Time dependent random variables : time+space
 - time: 1, 2, 3, 4, --. $+ \in (0 + \infty)$
 - space: Ω
 - state: $X(t,\omega)$ where $t \in Time, \omega \in \Omega$

Examples:



• Trajectory of a particle

Noise

• Rain

• Messages in a communication bus

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Examples:

• CPU usage

microcontrollers power consumption



2 Continuous time process

Markov process

only the most recent state counts

Stochastic process X(t) is Markov if for any $t_1 < ... < t_n < t$ and any sets $A; A_1, \ldots, A_n$



Markov chain



- discrete Markov process
- the state at time t+1 depends only on the state at time t

$$p_{ij}(t) = P \{ X(t+1) = j \mid X(t) = i \}$$

= $P \{ X(t+1) = j \mid X(t) = i, X(t-1) = h, X(t-2) = g, \ldots \}$

Transition probability:

$$p_{ij}^{(\mathbb{L})}(t) = \mathbf{P} \{ X(t+A) = j \mid X(t) = i \}$$

Homogenous Markov chain

• Transition pbb does not depend on the time

• Transition matrix

(p_{11}	p_{12}	• • •	p_{1n}	
	p_{21}	p_{22}	• • •	p_{2n}	
Į	÷	÷	÷	÷	
ĺ	p_{n1}	p_{n2}	•••	p_{nn}	J

Transition in 2 steps



$$p_{ij}^{(2)} = P\{X(2) = j \mid X(0) = i\}$$

$$= \sum_{k=1}^{n} P\{X(1) = k \mid X(0) = i\} \cdot P\{X(2) = j \mid X(1) = k\}$$
$$= \sum_{k=1}^{n} p_{ik} p_{kj} = (p_{i1}, \dots, p_{in}) \begin{pmatrix} p_{1j} \\ \vdots \\ p_{nj} \end{pmatrix}.$$

Transition pbb in two steps



Probabilities at time t

Transition matrix M of a homogenous chain



Description via a Transition diagram



Transition diagram

2 users: active user disconnects with pbb 0.5 inactive user connects with ppb 0.2 X= number of active users



Steady state distribution

"eventually it does not depend on the initial state"

A collection of limiting probabilities

$$\pi_x = \lim_{h \to \infty} P_h(x)$$

is called a steady-state distribution of a Markov chain X(t).

It is not clear in advance that a steady-state distribution exists

Another name used: stationary distribution

Example: no steady state distribution

random walk in a bipartite graph



proces periodychy

Computing steady state distribution

$$P_{h}P = P_{0}P^{h}P = P_{0}P^{h+1} = P_{h+1}.$$

$$(\Pi_{0},\Pi_{1}) \begin{pmatrix} 0 \\ 0 \end{pmatrix} \ddagger_{0} \Pi_{0} \\ 0 \end{pmatrix} \ddagger_{0} \Pi_{0}$$
this is a system of linear equations. Moreover:
$$\sum_{\nu} \Pi_{\nu} = 4$$

(the probabilities must sum up to 1)

Weather example cnt

$$(\pi_1, \pi_2) = (\pi_1, \pi_2) \begin{pmatrix} 0.7 & 0.3 \\ 0.4 & 0.6 \end{pmatrix} = (0.7\pi_1 + 0.4\pi_2, 0.3\pi_1 + 0.6\pi_2).$$

$$\begin{cases} 0.7\pi_1 + 0.4\pi_2 &= \pi_1 \\ 0.3\pi_1 + 0.6\pi_2 &= \pi_2 \end{cases} \Leftrightarrow \begin{cases} 0.4\pi_2 &= 0.3\pi_1 \\ 0.3\pi_1 &= 0.4\pi_2 \end{cases} \Leftrightarrow \pi_2 = \frac{3}{4}\pi_1.$$

$$\pi_1 + \pi_2 = \pi_1 + \frac{3}{4} \pi_1 = \frac{7}{4} \pi_1 = 1,$$

 $\underline{\pi_1 = 4/7}$ and $\underline{\pi_2 = 3/7}.$

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Existence of stationary distribution

A Markov chain is regular if

 $p_{ij}^{(h)} > 0$

for some h and all i, j. That is, for some h, matrix $P^{(h)}$ has only non-zero entries, and h-step transitions from any state to any state are possible.

Any regular Markov chain has a steady-state distribution.

Example: random walk on a bipartite graph --- this property does not hold!

h is even: H is odd:

Breaking periodicity

The trick is to modify the transition function T:

- with pbb 0.5 keep the old state
- with pbb 0.5 change state according to T

Algorithms based on Markov chains

Example: choose a maximal independent set in a graph at random (with uniform probability)

Difficult:

we can create a maximal independent set (e.g. via a greedy algorithm), enumerating all MIS might be extremely hard

Approach: via a Markov chain

- states = independent sets
- transitions: simple modifications (removing or adding nodes)
- ... so that the steady distribution is uniform

Uniform steady distribution

The case of double stochastic matrix: sum of each row is 1 (must be) sum of each column is 1 (not for all transition matrices)

Theorem:

stationary distribution is uniform for a double stochastic transition matrix

Proof: to check that:

$$\begin{pmatrix} 1 & \frac{1}{n} \\ \frac{1}{n} \end{pmatrix} \cdot \begin{bmatrix} p \\ p \end{bmatrix} = \begin{pmatrix} 2 \\ \frac{1}{n} \\ \frac{1}{n}$$

Checking one column:

$$\frac{1}{h} P_{1j} + \frac{1}{h} P_{2j} + \frac{1}{h} P_{nj} = \frac{1}{h} (P_{1j} + - + P_{nj}) = \frac{1}{h} (\frac{1}{h} - \frac{1}{h})$$

Uniform steady distribution – special case: Symmetric transition matrix



Obviously: the matrix is double stochastic → the steady distribution is uniform!

Absorbing states

No exit from an absorbing state



Absorbing state example



Algorithms based on absorbing state

Random walk based on a Markov chain

Run simulation, eventually you are trapped in an absorbing state = a good state, where there is nothing to improve

Counting processes

e.g. Bernoulli trials State in time t = number of succeses in steps 1 through t

Expected number: p't

Counting processes

state is a counter counter is nondecreasing

Examples:

- the number of incoming cars on a bridge
- the number of emails arrived

Binomial proces

independent Bernoulli trials

counter= number of successes time frame Δ : one Bernoulli trial per Δ seconds

Expected number of successes:

$$\mathbf{E}\left\{X\left(\frac{t}{\Delta}\right)\right\} = \frac{t}{\Delta}p$$

Expected number of successes per second (arrival rate):

$$\lambda = \frac{p}{\Delta}$$

Bernoulli counting process - interarrival time

 $\mathcal{Y} = \frac{\mathcal{V}}{\mathcal{S}}$

 $T = Y\Delta$

where Y has geometric distribution

$$\begin{split} \mathbf{E}(T) &= \mathbf{E}(Y)\Delta = \frac{1}{p}\Delta = \frac{1}{\lambda};\\ \mathrm{Var}(T) &= \mathrm{Var}(Y)\Delta^2 = (1-p)\left(\frac{\Delta}{p}\right)^2 \text{ or } \frac{1-p}{\lambda^2}. \end{split}$$

Continuous counting process

a limit of Bernoulli counting process with time frame $\Delta \rightarrow 0$

The number of frames during time t increases to infinity,

$$n = \frac{t}{\Delta} \uparrow \infty \text{ as } \Delta \downarrow 0.$$

The probability of an arrival during each frame is proportional to Δ , so it also decreases to 0,

 $p = \lambda \Delta \downarrow 0$ as $\Delta \downarrow 0$.

Continuous counting process:

Poisson as a limit of Binomial

Then, the number of arrivals during time t is a Binomial(n, p) variable with expectation

$$\mathbf{E}X(t) = np = \frac{tp}{\Delta} = \lambda t.$$

$$X(t) = \text{Binomial}(n, p) \rightarrow \text{Poisson}(\lambda)$$

Based on the following theorem:

$$\lim_{\substack{n \to \infty \\ p \to 0 \\ np \to \lambda}} {n \choose x} p^x (1-p)^{n-x} = e^{-\lambda} \frac{\lambda^x}{x!}$$

Continuous counting process

The interarrival time T becomes a random variable with the c.d.f.

$$F_T(t) = P\{T \le t\} = P\{Y \le n\}$$
$$= 1 - (1 - p)^n$$
$$= 1 - \left(1 - \frac{\lambda t}{n}\right)^n$$
$$\to 1 - \bar{e}^{\lambda t}.$$

because $T = Y\Delta$ and $t = n\Delta$ Geometric distribution of Y

because $p = \lambda \Delta = \lambda t/n$ This is the "Euler limit": $(1 + x/n)^n \to e^x$ as $n \to \infty$

The interarrival time converges to exponential distribution

Continuous counting process

let us inspect the time for k arrivals

 $P\{T_k \leq t\} = P\{k \text{-th arrival before time } t\} = P\{X(t) \geq k\}$ where T_k is $\text{Gamma}(k, \lambda)$ and X(t) is $\text{Poisson}(\lambda t)$.

Applications

What is the probability that in time T more than k requests arrive for a webpage P?

We assume that λ is known (λ requests per minute)

Continuous counting process:

Poisson as a limit of Binomial

However: if we take a time interval that is not very small with respect to λ , then some differences occur:

Binomial: at most one arrival in the interval

Poisson: more than one possible

... even if the arrival rate is the same λ , there is a greater variance for Poisson

Solution for λ =7 hits per minute , assumed Poisson process

Pbb for 10000 hits within 24 hours?

<u>Solution</u>. The time of the 10,000-th hit T_k has Gamma distribution with parameters k = 10,000 and $\lambda = 7 \text{ min}^{-1}$. Then, the expected time of the k-th hit is

$$\mu = \mathbf{E}(T_k) = \frac{k}{\lambda} = \underline{1,428.6 \text{ min or } 23.81 \text{ hrs.}}$$

$$\sigma = \operatorname{Std}(T_k) = \frac{\sqrt{k}}{\lambda} = 14.3 \text{ min.}$$

Pbb of more than 10000 hits?

A shortcut: CLT (we do not have to care about Gamma distribution!)

$$P\left\{T_k < 1440\right\} = P\left\{\frac{T_k - \mu}{\sigma} < \frac{1440 - 1428.6}{14.3}\right\} = P\left\{Z < 0.80\right\} = \underline{0.7881}.$$

Conclusion

Easy way to solve many problems regarding required capacity ...